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TITLE OF INVENTION

METHOD AND SYSTEM FOR DYNAMIC CONFIGURATION OF ACTIVATORS IN A CLIENT-SERVER ENVIRONMENT

APPLICANT(S) FOR DO/EO/US

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Applicant herewith submits to the United States Designated/ Elected Office (DO/EO/US) the following items under 35 U.S.C. 371:

1. ☒ This is a **FIRST** submission of items concerning a filing under 35 U.S.C. 371.
2. ☐ This is a **SECOND** or **SUBSEQUENT** submission of items concerning a filing under 35 U.S.C. 371.
3. ☒ This is an express request to begin national examination procedures (35 U.S.C. 371(f)) at any time rather than delay examination until the expiration of the applicable time limit set in 35 U.S.C. 371(b) and PCT Articles 22 and 39(1).
4. ☒ A proper Demand for International Preliminary Examination was made by the 19th month from the earliest claimed priority date.
5. ☒ A copy of the International Application as filed (35 U.S.C. 371(c)(2))
 - a. ☐ is transmitted herewith (required only if not transmitted by the international Bureau).
 - b. ☒ has been transmitted by the International Bureau.
 - c. ☐ is not required, as the application was filed in the United States Receiving Office (RO/US)
6. ☐ A translation of the International Application into English (35 U.S.C. 371(c)(2)).
7. ☐ Amendments to the claims of the International Application under PCT Article 19 (35 U.S.C. 371(c)(3))
 - a. ☐ are transmitted herewith (required only if not transmitted by the International Bureau).
 - b. ☐ have been transmitted by the International Bureaus.
 - c. ☐ have not been made; however, the time limit for making such amendments has NOT expired.
 - d. ☐ have not been made and will not be made.
8. ☐ A translation of the amendments to the claims under PCT Article 19 (35 U.S.C. 37(c)(3)).
9. ☐ An oath or declaration of the inventor(s) (35 U.S.C. 371(c)(4)).
10. ☐ A translation of the annexes to the International Preliminary Examination Report under PCT Article 36 (35 U.S.C. 371(c)(5)).

Items 11. to 16. below concern document(s) or information included:

11. ☒ An Information Disclosure Statement under 37 CFR 1.97 and 1.98.
12. ☐ An assignment document for recording. A separate cover sheet in compliance with 37 CFR 3.28 and 3.31 is included.
13. ☐ A **FIRST** preliminary amendment.
☐ A **SECOND** or **SUBSEQUENT** preliminary amendment.
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15. ☐ A change of power of attorney and/or address letter.
16. ☒ Other items or information:

INTERNATIONAL SEARCH REPORT

- 17.
- ☒
- The U.S. National Fee (35 U.S.C. 371(c)(1)) and other fees as follows:

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METHOD AND SYSTEM FOR DYNAMIC CONFIGURATION OF ACTIVATORS IN A CLIENT-SERVER ENVIRONMENT

FIELD OF THE INVENTION

The present invention is directed to configuring a server computer in a client-server environment and, in particular, to dynamically configuring activators.

5 BACKGROUND OF THE INVENTION

With the rise of the interconnected computer networks such as the Internet, it is possible to construct complex transaction-based applications that are distributed over several networked computers. In the simplest scenario, in general, these transaction-based applications function in the following way. A software application program, which
10 executes on a client, initiates a transaction that requires access to services provided by a distant computer, called a server. Examples of these services could be an update to a database such as a bank's database, an execution of a purchase order such as in the case of purchase of a security and the like. Typically, the client sends a "request" message to the server, which then sends a "response" message containing a response to the request.

15 Typically, the server is not a single computer, rather a collection of interconnected heterogenous computers. The request message is then be formatted in such a way that all the interconnected computers can understand and respond to the request message. If the collection of interconnected computers is configured in an object-oriented programming model, then software object (or objects) that are capable of working together
20 to provide a response to the request message can be distributed among the several computers. But in order to access the objects from a remote computer the objects must somehow publish their existence, their addresses, their properties, the services they provide, and other details to the "outside" world. Then, a client may be able to use the services provided by sending a request message in a manner similar to making a remote procedure
25 call ("rpc") and obtaining a response to that message.

Various paradigms exist as a result of the need to standardize the methods by which objects can be distributed and accessed over a network. These are Microsoft Corporation's Distributed Component Object Model (DCOM), JavaSoft's Java/Remote

Method Invocation (Java/RMI), and Object Management Group's Common Object Request Broker Architecture (CORBA).

Though some differences are present among these models, they principally work in the following way. Objects that provide services are typically located on servers.

5 These objects are queried by applications running on clients using a specified data communication transport layer protocol—the Object Remote Procedure Call (ORPC) for DCOM; the Java Remote Method Protocol (JRMP) for Java/RMI; and the Internet Inter-ORB Protocol (IIOP) for CORBA. A client suitably formats a query message in the appropriate protocol language and transmits the query message, which is routed to the
10 appropriate server, whereupon it is executed, and a response message is formatted and routed back to the client. As referred to herein, the term “object” may mean the object definition, associated operations, attributes, etc., and implementation for that object. As will be appreciated by those of skill in the art, at times the term “object type” is used to refer to the definition of the operations and attributes that software external to the object may use
15 to examine and operate upon the object. The “object type” is also known as the “interface.” Also, the term “object” may be used to refer to an actual run-time instance of an object and will be made clear by the context.

A server configured to be a Java/RMI server comprises objects that have predefined interfaces, which can be used to access the server objects remotely from another
20 machine's Java Virtual Machine (JVM). A Java/RMI server object interfaces declare a set of methods that indicate the services offered by that server object. A program resident on the server called an RMI Registry stores and makes available to clients information about server objects. Typically, a client object obtains information regarding the methods and other properties of a server object by performing an operation such as “lookup” for a server
25 object reference. This lookup typically works by the client object specifying an address in the form of a Universal Resource Locator (URL) and transmitting the address to the server's RMI Registry.

The clients and servers also include interceptors. The interceptors provide hooks to programmers to execute their piece of code at certain points during ORB. Typical
30 uses of the interceptors include: transaction service integration, security message compression and encryption, fault tolerance and other operations such as tracing, profiling, debugging, logging.

In CORBA, each CORBA object transparently interacts with an Object Request Broker (ORB), which provides a means to access either local or remote objects. The ORB is essentially a remote method invocation facility, and forms the lowest layer of the several layers in CORBA. Each CORBA server object exposes a set of methods, and it declares its interface. A CORBA client obtains an object reference and determines which methods are provided by the object. A CORBA client needs only two pieces of information: a remote object's name, and how to use its interface. The ORB is responsible to locate the object, provide a vehicle by means of which a query is transmitted to a server object and a response is transmitted back to the client object. In general, a CORBA object interacts with an ORB by either using an ORB's interface or using an Object Adapter.

There are two kinds of object adapters, the Basic Object Adapter (BOA) and the Portable Object Adapter (POA). The BOA (or the POA) typically has methods for activating and deactivating objects, and for activating and deactivating the entire server. These are intended for systems where the ORB and the server are separate programs or even on separate machines. Different vendors supplying CORBA-compliant servers ordinarily choose one or the other of these methods of an object-ORB interaction.

As described above, CORBA objects take form within server applications. In a server, CORBA objects are implemented and represented by programming language functions and data. The programming language entities that implement and represent CORBA objects are called servants. A servant is an entity that provides a body for a CORBA object, and for this reason, the servant is said to incarnate the CORBA object.

Object adapters such as the CORBA-standard Portable Object Adapter (POA) mediate between an ORB and a set of programming language servants. In general, though there could be many instances of POAs to support CORBA objects of different styles, and all server applications have at least one POA called the *Root* POA. Each POA instance represents a grouping of objects that have similar characteristics. These characteristics are controlled via POA policies that are specified when a POA is created. The Root POA, which is present in all server applications, has a standard set of policies. POA policies are a set of objects that are used to define the characteristics of a POA and the objects created within it. The CORBA standard specifies that interfaces for POA, POA manager (which is a class to manage multiple POAs) and the POA policies should be defined in a standard module.

The above discussed technologies have been utilized in the Internet. However, the next phase of the Internet revolution is predicted to be interconnection of isolated Internet systems with the systems that run the business to create a responsive, flexible, scalable, and differentiated eBusiness enterprise. The information systems that connect eBusiness with the enterprise are coming to be known as *enterprise portals*.

Enterprise portals can act as new storefronts, new front-offices, new sales and support agents for the enterprise, with profound implications. Enterprise portals will leverage an enterprise's processes, data, and transactions on the Internet. They will simplify access to a mix of Internet and non-Internet applications, built with heterogeneous formats, platforms, protocols and software. This universe of diverse content, interactions, transactions, and application functions will require new methods of organization and management, particularly to keep up with frequent changes in the business.

The enterprise portals are becoming popular because the Internet affords large opportunities to extend business and invent new business models. Enterprise portals hold the key for these established companies to "bricks and clicks" strategies that weave online services with their existing channels, production facilities, and other business elements into a powerful combination.

Enterprise portals require scalable, flexible and open distributed platforms, the ability to leverage and extend the full range of resources inside and outside of the corporate walls, and the ability to accommodate rapid, continual change. At their root, enterprise portals host a different kind of application -- the *composite application*.

Composite applications combine new business logic with existing logic, processes, and data to meet business needs. A new system is required to provide all of the components of an enterprise portal infrastructure. This new set of server-based products marshals the company's expertise in design and support of enterprise systems from diverse applications using object and component technologies.

SUMMARY OF THE INVENTION

Therefore, the present invention provides a flexible, robust, CORBA 3.0 middleware engine. Adaptive Runtime Technology™ (ART) framework, a microkernel architecture that supports dynamic configurations of the architecture. It should be noted that the present invention is described below by discussing various ports of the ART framework.

In particular, the present invention provides a computer implemented method of activating a process. The method includes the steps of generating one or more first plug-ins each configured to activate a target process, dynamically registering the first plug-ins with a second plug-in and permanently storing information relating to each registered first plug-in.

The method can further include the steps of storing a flag for each registered first plug-in, perpetually activating the corresponding target process if the flag is set to a first state, and activating the corresponding target process upon a request if the flag is set to a second state.

The method can also include the steps of generating an exception to indicate that a target process is inactive when its flag is not set to the first state or the second state, providing a unique identifier for each target process, and sending and receiving a message between the first and second plug-ins using the identifiers. The message can include information relating to a state change of the target processes, and the state includes an activated state and a deactivated state.

The present invention further provides a server computer in a client-server computer system. The system includes one or more first plug-ins each configured to activate a target process, and a second plug-in configured to dynamically register the first plug-ins and to permanently store information relating to the registered first plug-ins.

The second plug-in of the server can also include a memory configured to store a flag for each registered first plug-in. The second plug-in can be further configured to perpetually activate target processes having their flags set at a first state and to activate target processes, upon receiving a request, having their flags set at a second state.

Moreover, the second plug-in can be further configured to generate an exception to indicate that the target process is inactive when the flag is not set to the first state or the second state.

The server computer can also include a first computer program object configured to provide a unique identifier for each target process and configured to send a message using the identifiers. The message includes information relating a state change of the target processes, and the state includes an activated state and a deactivated state.

BRIEF DESCRIPTION OF THE DRAWINGS

Preferred features of the present invention are disclosed in the accompanying drawings, wherein similar reference characters denote similar elements throughout the several views, and wherein:

FIG. 1 is a diagram illustrating a computer network for the distributed
5 objects of the present invention;

FIG. 2 is a block diagram of a typical computer of the present invention;

FIG. 3 is a diagram illustrating communication among a locator, activator
and target process;

FIG. 4 is a diagram illustrating class association between a location plug-in
10 and an activation plug-in;

FIG. 5 is a diagram illustrating a class hierarch of an object of the present
invention;

FIG. 6 is a diagram illustrating a class hierarch of an *IT_ActivatorPlugIn*
class;

FIG. 7 is a diagram illustrating a class hierarch of an *IT_ActivatorImpl* class;

FIG. 8 is a diagram illustrating a class hierarch of an
15 *IT_ProcessTerminationMonitor* class;

FIG. 9 is a diagram illustrating a class hierarch of an *IT_ProcessImpl* class;

FIG. 10 is a diagram illustrating a class hierarch of an *IT_ProcessLauncher*
20 class;

FIG. 11 is a diagram illustrating a class hierarch of internal tables of the
present invention;

FIG. 12 is a diagram illustrating IR management interfaces;

FIG. 13 is a diagram illustrating a class hierarch of OA independent
25 information;

FIG. 14 is a diagram illustrating three scenarios of moving a POA from one
ORB instance to another; and

FIG. 15 is a diagram illustrating a generic sever model.

DETAILED DESCRIPTION OF THE INVENTION

Referring to FIG. 1, distributed objects of the present invention are located on one or more computers linked together by a computer network exemplified in a network 10. In particular, the network 10 includes a computer 12 coupled to a network 14. The network 14 can further include a server, router or the like 16 in addition to other computers 18, 20, and 22 such that data, instructions and/or messages can be passed among the networked computers. A mass storage devices 24 may also be connected to the server 16 or to any of the computers. Further, some computers 12, 18 may include an independent network connection between them, whereas other computers 20, 22 may not include such a connection. Various ways to design, construct and implement the computer network as known in the art are contemplated within this invention.

Referring to FIG. 2, each computers 12, 16, 18, 20, and 22 includes a processing unit 42, a primary storage device 44 and a secondary storage device 46. The processing unit 42 can be, but not limited to, a central processing unit (CPU), or multiple processors including parallel processors or distributed processors. The primary memory device 44 includes random access memory (RAM) and read only memory (ROM). The RAM stores programming instructions and data, including distributed objects and their associated data and instructions, for processes currently operating on the processor 42. The ROM stores basic operating instructions, data and objects used by the computer to perform its functions. The secondary storage device 46, such as a hard disk, CD ROM, magneto-optical (optical) drive, tape drive or the like, is coupled bidirectionally with processor 42. The secondary storage device 46 generally includes additional programming instructions, data and objects that typically are not in active use by the processor, although the address space may be accessed by the processor, *e.g.*, for virtual memory or the like.

Furthermore, each of the above described computers can include an input/output source 50 that typically includes input media such as a keyboard, pointer devices (*e.g.*, a mouse or stylus) and the like. Each computer can also include a network connection 52. Other variations of the above discussed the computer and its components available to one of skill in the art are also contemplated within the present invention.

In the present invention computer network is defined to include a set of communications channels interconnecting a set of computer systems that can communicate with each other. The communications channels can include transmission media such as, but not limited to, twisted pair wires, coaxial cable, optical fibers, satellite links, and/or digital microwave radio. The computer systems can be distributed over large, or "wide" areas (*e.g.*,

over tens, hundreds, or thousands of miles. WAN), or local area networks (e.g., over several feet to hundreds of feet, LAN). Furthermore, various local- and wide-area networks can be combined to form aggregate networks of computer systems. One example of such a network of computers is the "Internet".

5 Having discussed a client-sever computing environment, the following is the ART framework briefly described above.

I. ART ACTIVATOR

10 For ART, activation functionality is preferably implemented as a plug-in to a generic ART daemon. The activation plug-in is responsible for activating a target process and has no persistent data to manage. A target process is an executable process containing one or more target endpoints.

15 One or more activation plug-ins may be configured within a location domain. There may be zero to many activation plug-ins on a single host. But an ART daemon with the activation plug-in preferably lives on the host where the target process resides. Also note that different location domains and their respective location plug-ins may share a single activation plug-in.

20 For any communications to occur in a location domain between the location plug-in and any activation plug-ins, each activation plug-in dynamically registers itself with the location plug-in via an *ActivatorRegistry* interface which is obtained at configuration time. The location plug-in is required to store every registered *Activator* in its *Activator List* in permanent storage. If the host where the location plug-in resides crashes, then it can be re-started without requiring the activation plug-ins to re-register or the location plug-in to obtain state information on all its IMR entries.

25 When the location plug-in gets control and there is an endpoint name match but the associated target process is not activated, the location plug-in checks the start-up flags in the static IMR entry for the associated target process. If the start-up flags are not set to "always" or "on demand", a *ProcessNotActive* exception is required to be returned to the client. Otherwise, the location plug-in first creates a target process dynamic IMR entry with a target process state of *STARTING*. In order to find an appropriate activation plug-in to
30 start a target process, the location plug-in will map the target endpoint's associated target process static IMR entry in the IMR to the registered *Activator* information. Finally, the

location plug-in creates a *ProcessMonitor* for notification of any target process state changes and sends a request to start the target process.

When the activation plug-in receives a request to start a target process, it identifies any invalid parameters and security violations. The activation plug-in then creates
5 a *Process* representing the target process and stores the associated *ProcessMonitor* in it. The activation plug-in returns the *Process* to the location plug-in. At this time, the location plug-in stores the *Process* in the associated target process dynamic IMR entry. Depending on the wait options of the start process request, the target process may or may not have been started.

10 At any point in time, the activation plug-in may return information about a target process state change. If the target process is up and running, the activation plug-in announces to the location plug-in that the target process was activated via the *ProcessMonitor*. The location plug-in updates the appropriate dynamic IMR entry and transitions the target process state to *RUNNING*. The activation plug-in then waits on the
15 target process. This allows target processes to be monitored. In addition, network traffic is reduced if the location plug-in has the target process's current state.

If the target process dies, the activation plug-in announces this state change to the location plug-in via the *ProcessMonitor*. The location plug-in deletes the dynamic IMR entries associated with the target process name in the IMR and sends a request to
20 remove the associated *Process*. For target process transparent restart, the location plug-in then checks the target process's start-up flags. If "disable start", then target process restart will not occur. If "on demand", target process restart occurs when a client request arrives for that target process again. If "always", then the location plug-in immediately attempts to restart the target process. If "on demand" or "always" and unsuccessfully started, an
25 exception is required to be returned to the client.

This section describes a low-level common OS-independent abstraction for remote process activation without compromising system security. A process is launched with a specified identity, running a specified program with specified arguments and environment. There are also capabilities to ask or force the process to terminate, and notice
30 when it does terminate. Defining the Activator at as low a level of abstraction as possible preserves the full flexibility of the underlying operating systems. More responsibility is given this way to the Locator, where the low-level abstraction can be used as appropriate.

1. Module IT_LocationCommon

This module includes common exceptions and data types amongst all Location-related interfaces.

```

5  module IT_LocationCommon
   {
       struct EnvVariable
       {
           string          name;
10          string          val;
       };
       typedef sequence<EnvVariable> EnvVariableList;
       typedef sequence<string> ArgumentList;
       struct StartUpInformation
15      {
           string          pathname;
           ArgumentList    arguments;
           EnvVariableList env_variables;
           string          user; // Real user ID of target
20      process
           string          group;// Real group ID of target
       process
           } ;

25      struct SecurityId
       {
           string principal;
       } ;
       typedef sequence<SecurityId> SecurityIdList;
30  } ;

```

2. Module IT Activation

This module includes four interfaces; `Process`, `ProcessMonitor`, `Activator`, `ActivatorRegistry`.

(1) The Process

The Activation plug-in provides an implementation of this interface. All *Process* object instances are transient. A *Process* is created by an *Activator* to represent a single target process that is to be started. The *Process* is returned to a *Locator* and stored in the target process's dynamic IMR entry in the IMR. The purpose of a *Process* is to provide a unique identifier for the target process for communications between the *Locator* and *Activator* since different *Locators* representing different location domains may share an *Activator* and there may be clashes with target process names.

If the entire host where the activation plug-in and all target processes reside goes down, then all *Process* references are invalid in the IMR but all target processes on that host have died also. Once the activation plug-in comes back up, it will re-register with one or more *Locators* and begin to receive requests to re-start any target processes which have a start-up option of *always*.

In the rare case when only the generic daemon with an activation plug-in goes down and not the entire host, then all *Process* references are invalid in the IMR. Once the *Activator* comes up it cannot monitor those target processes that it had previously created and started but they may still be running on that host. In addition, a *Locator* may send a request to start a target process that is already running once the *Activator* re-registers.

During a *kill()* operation, the associated running target process on the system is terminated if the proper access rights are allowed to do so. During a *remove 0* operation, only the specified *Process* will be removed. Note that if the target process associated with this *Process* is still up and running, a *ProcessStillActive* exception is returned.

A management interface for the *Locator* will be provided to start and stop target processes. In this case, the *Locator* acts as an intermediary for a client request to start/stop processes. The *Locator* extracts credentials from the client request and then sets its own credentials to that of the client before making the request to the *Activator*.

The *Process* object maintains three timestamps. The startup timestamp is the date/time the target process is started. The died timestamp is the time/date the target process has died. The expiration timestamp is the time/date to remove the *Process* object after the target process has died. It should be noted that the *Activator* does not remove *Process* objects but the *IT_Activation::Activator::create-process()* requestor does. The *Activator* can prevent *Process* objects from dangling. For example, if the *Locator* goes down and an *Activator* never receives a *remove()* request. An expiration timestamp (default

24 hours, a configurable) on Process objects is obtained from configuration. If the Locator does not remove the Process, then the Activator will once the expiration time expires.

```

module IT_Activation
5  {
    interface Process
    {
        // Exceptions
        //
10     exception ProcessStillActive {};
        exception ProcessNotExist {};
        exception InvalidSignal {};
        // This ProcessMonitor is created by the Locator for
notification
15     // of any target process state changes. In the rare
event that
        // the generic daemon with the Location plug-in goes
down, then
        // all ProcessMonitor references no longer exist for all
20     active
        // target processes in the IMR. However when the
Locator comes
        // back up, it can recover. For each and every active
target
25     // process, a new ProcessMonitor is created and the
attribute
        // in its corresponding Process is reset.
        attribute ProcessMonitor
process_monitor;
30     // The start-up time stamp is used for informational
purposes.
        //
        readonly attribute long startup timestamp;
35     // Allows the Activator to garbage collect a Process
that is not
        // removed but the associated target process is dead.

```

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```

//
// The died timestamp is the time/date the target
process has died.
// The expiration timestamp is the time/date to remove
5 the Process
// object after the target process has died.
//
readonly attribute long          died_timestamp;
readonly attribute long          expiration_timestamp;
10
// The state attribute is used by a Locator to determine
the
// cause of a failed target process.
//
15
enum ProcessStateValue
{
    CREATING_PROCESS,
    CREATE_PROCESS_FAILED,
20    STARTING_PROCESS,
    START_PROCESS_FAILED,
    DIED,
    RUNNING
} ;
25 readonly attribute ProcessStateValue          state ;

// The pid attribute is required for killing processes,
attaching
// processes to debuggers, etc.
30
//
typedef unsigned long          ProcessId ;
readonly attribute Process ID          pid ;
// If an announcement that the process was not started
or died,
35 // remove the Process.
//

```


Locator can obtain the *state* attribute in the *Process* and return the appropriate exception to the client.

The activation plug-in also helps in monitoring target processes by waiting for its children to die that it had exec-ed. When a target process has died, the *Activator* will set the *state* to *DIED* in the *Process* and announce its death to the Locator. The location plug-in then deletes the dynamic IMR entries associated with this target process name, deletes the associated *ProcessMonitor* and sends a request back to remove the *Process*.

The exceptions on the *announce_process()* operation are required to be logged. For *ProcessNotExist*, the target process that was started can be killed. This indicates that a target process was started by an *Activator* but there is no target process static IMR entry in the Implementation Repository.

In the rare case when the generic daemon with location plug-in goes down, then all *ProcessMonitor* references no longer exist for all active target processes in the IMR. However when the Locator comes back up, it can recover. For each and every active target process, a new *ProcessMonitor* is created and the attribute in its corresponding *Process* is reset.

```

module IT_Activation
{
20     interface ProcessMonitor
        {
            // Exceptions raised with this target process instance
and
            // information found in the IMR.
25         //
            exception ProcessNotExist { } ;
            exception ProcessNotActive { } ;
            exception ProcessAlreadyActive { } ;

30         enum ProcessEventValue
            {
                STARTED,           // Target process was
started
                NOT_STARTED,       // Target process was not
35 started

```

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```

        DIED                                // Target process has died
    } ;

    // Announces a target process event to the location
5  plug-in
    //

    void
    announce_process (
10         in ProcessEventValue      event_type
    ) raises (
        ProcessNotExist,
        ProcessNotActive,
        ProcessAlreadyActive
15     ) ;
    } ; // End of Process Monitor Interface
} ;

```

(3) The Activator

20 The Activation plug-in provides an implementation of this interface. All *Activator* object instances are transient. The purpose of the *Activator* is to identify itself to one or more location plug-ins for target process activation to occur. There is only one *Activator* per activation plug-in. But note that the *Activator* may be registered with one or more Locators with the same name or a different name.

25 For any communications to occur in a location domain between the location plug-in and any activation plug-ins, each activation plug-in preferably registers itself with the location plug-in. Once registered, the *Activator* is stored on disk. The *Activator List* is a list of registered *Activator(s)* with a location plug-in for the location domain. The *Activator List* is managed by the location plug-in and resides in permanent storage. An
30 entry in the *Activator list* contains an activator name, an *Activator* reference, and the owner of the entry. The activator name is unique within a location domain. If several Locators share an *Activator*, it may or may not be registered with the same name.

During a *create_process()* operation, the *Activator* performs validations on the input parameters sent. If these parameters are not valid, an *InvalidConfiguration*
35 exception is returned to the client. If the path name of the executable process does not exist,

an *ExecutableNotFound* exception is returned to the client. In addition, the next two paragraphs identify security violations.

The *Activator* preferably ensures that the security identity activating the target process has the proper access rights to do so. This security identity identifies the real user and real group IDs of a target process. Since all target processes exec-ed from the activation plug-in inherit the activation plug-in's real user and real group IDs, it is the responsibility of the *Activator* to ensure that the security identity configured in the target process's static IMR entry maps to the real user and real group ID of the activation plug-in. If not, an *CORBA::NO_PERMISSION()* exception is required to be sent back to the client.

In addition on UNIX, the activation plug-in preferably determines if the effective user and group IDs may have been given additional permissions via set-uid and/or set-gid bits on the ART daemon with the activation plug-in itself or the target process executable file. In this case, the real user and group IDs may map to the user and group IDs in the IMR but the effective user and group IDs may have been given additional permissions not intended by the entity that established the static target process IMR entry. If so, an *CORBA::NO_PERMISSION()* exception is required to be sent back to the client.

Once the above validations are successful, then the activation plug-in creates a *Process* representing the target process and stores the *ProcessMonitor* in it. If the *create_process()* operation is specified with the *NO_WAIT* option, the activation plug-in immediately returns the *Process* to the location plug-in and adds the *Process* to its in-memory *Process* list. For the *NO_WAIT* option, this is done before actually creating and starting the target process on the host. In this case, the notification of target process activation is via the *ProcessMonitor*. If activation is not successful, the activation plug-in implicitly removes the *Process* object.

If the *create_process()* operation is specified with the *WAIT* option, then the Locator waits for the *Activator* to actually create and start the specified target process on the host. If successful, the activation plug-in returns the *Process* representing the target process to the location plug-in and adds the *Process* to its in-memory *Process* list. If not, a *CreateProcessFailed* or *StartProcessFailed* exception is raised and a null value is returned.

If the *WAIT* option is chosen, the *IT_Activaiton::ProcessMonitor::announce process(STARTED/NOT_STARTED)* are not called as the *create_process()* operation returns successfully or raise an exception.

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The *WAIT/NO_WAIT* option on the *create_process0* operation is determined by the Locator. During a client request invocation, the Locator will use the *WAIT* option to start a target process. However at Locator start-up time, the Locator may make several requests to several Activators to start target processes with the *NO_WAIT* option.

If there is a rate-limit specified in the static IMR entry for the target process, the Locator may invoke *create_process()* operation one or more times. For each unsuccessful startup attempt, the activation plug-in implicitly removes the *Process* object. If the number of attempts exceeds the maximum number of retries configured in a target process static IMR entry, the troublesome exception is required to be returned to the client.

```

module IT_Activation
{
    interface Activator
    {
        exception ExecutableNotFound { } ;
        exception CreateProcessFailed { } ;
        exception StartProcessFailed { } ;
        exception InvalidConfiguration { } ;

        // Activator attributes
        //
        readonly attribute string host;
        readonly attribute string user;
        readonly attribute string group;

        // Creates and returns an object reference representing
        the process.
        // This operation is also used to retry a failing
        target
        // process. It is the responsibility of the caller to
        // remove the Process when no longer required.
        //

        Process
        create_process (

```



```

        in IT_LocationCommon : : StartUpInformation startup
information,
        in ProcessMonitor          callback,
        in booleans                 wait_for_completion
5          ) raises (
            ExecutableNotFound,
            InvalidConfiguration,
            CreateProcessFailed,
            StartProcessFailed,
10          ) ;
    }; // End of Activator Interface
} ;

```

(4) The Activator Registry

Referring to FIGs. 3 and 4, the Location plug-in provides an implementation of this interface. Preferably, there is only one *ActivatorRegistry* per Locator. Also note that one or more *ActivatorRegistries* may be available to an activation plug-in if several Locators share an activation plug-in. The purpose of the *ActivatorRegistry* is to provide a means for an activation plug-in to dynamically register with a Locator.

When a generic daemon with the activation plug-in is first started, it will obtain an object reference to this interface via some configuration mechanism. The activation plug-in will create an *Activator* and register it with the location plug-in. Once an *Activator* is registered, the location plug-in will search the static target process IMR entries in the IMR for those target processes configured with the same Activator name and a start-up option of *always*, which indicates that the target process is required to be up and running at all times. If a target process is determined to be no longer active, then the locator *implicitly* sends *create_process()* request to the activation plug-in. As a result at Activator registration time, there may be several requests to start target processes. This provides support for those instances when the host where the activation plug-in resides goes down.

The following actions provide support for the rare case when only the generic daemon with an activation plug-in goes down and not the entire host. If a generic daemon with an activation plug-in is *not* gracefully shutdown, its *Activator* may still be registered with a location plug-in. If the location plug-in attempts to use an *Activator* and receives an OBJECT_NOT_EXIST, the location plug-in will *implicitly* remove the

Activator from the Activator List. However, if the activation plug-in is restarted and attempts to register with the location plug-in again and there is a matching *Activator* name attribute in the Activator List, the location plug-in can implicitly try to use the already registered *Activator*. If the location plug-in receives an `OBJECT_NOT_EXIST`, it will

5 implicitly remove this *Activator* and allow registration of the new *Activator*. Otherwise, the activation plug-in will receive an *ActivatorAlreadyRegistered* exception.

```

module IT_Activation
{
10     interface ActivatorRegistry
        {
            // The Entry is the unit of information managed by the
            registry for a
            // given Activator instance. The "name" is the
15 administrative name used
            // for the Activator. It is used by the location
            daemon to decide which of
            // multiple Activator instances running on a particular
            host machine to be
20 // used to start a given server process. The "owner"
            is the security ID
            // of the Activator itself. This value is used to
            ensure that only
            // administrators or the Activator itself may modify
25 specific information
            // related to the Activator instance.
            //

            struct Entry
30     {
                string                name;
                Activator              activator_ref;
                IT_LocationCommon : : SecurityId    owner;
            } ;
35     typedef sequence <Entry> EntryList;
    exception ActivatorAlreadyRegistered {

```

```
        Entry activator_entry;
    } ;
    exception ActivatorNotRegistered {
        string name;
5        };
        // Register the existence of the indicated Activator
instance with the
        // Location daemon.
        //
10
        void
        add (
            in Entry activator_entry
        ) raises (
15            ActivatorAlreadyRegistered
        ) ;
        // Unregister the indicated Activator instance with the
location
        // daemon.
20        //

        void
        remove (
            in string name
25        ) raises (
            ActivatorNotRegistered
        ) ;
        // Return the information about the Activator instance
registered with
30        // the Location daemon specified by the provided "name"
        //

        boolean
        find (
35            in string name,
            out Entry activator_entry
        ) ;
```

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// Return a list of the Activator instances registered
with the location

// daemon.

//

EntryList

list () ;

} ; // End of Activator Registry Interface

} ;

3. Other aspects of Art Activator

In the rare case when only the generic daemon with an activation plug-in goes down and not the entire host, then all *Process* references are invalid in the IMR. Once the Activator comes up it cannot monitor those target processes that it had previously created and started but they may still be running on that host. In addition, a Locator may send a request to start a target process that is already running once the Activator re-registers.

In order to solve the above described problem, the present invention provides:

- The Activation plug-in keeps persistent information about the target processes it has started.
- *Process* is created as a persistent rather than a transient object.
- The Locator gets a POACallback via an announcement that a POA with a persistent lifespan policy was created. The Locator can determine if a target endpoint is still active for the associated target process via *is_active()* operation on the POACallback. If one target endpoint is active, then its associated target process is still active. Then the issue is, which target endpoint does the Locator choose if no target endpoints for a target process were configured for a heartbeat purpose. A configuration requirement can be made such that one target endpoint is configured for the heartbeat feature per target process.

In order to ensure that a Locator is talking to an Activator, the present invention provides:

- The Activator can be made only to accept requests to create a process from a locator that the activator knows about from its own configuration. So, the activator's configuration can have a list of locator IORs to register with. This IOR can have a TAG_SEC_NAME containing the Security Name of the locator. The activator verifies the Security Name of the client every time *create_process()* is called. (The Security Name can be obtained from *SecurityCurrent.get_attributes()*)

Providing dynamic registration and/or registering Activators via some Configuration mechanism is contemplated within the present invention. This will determine persistent or transient object for Activator registration.

A management interface to register Activators with Locators can also be provided in the present invention.

Furthermore, *IT Activation::Process::ProcessId* can be a string making it more portable across various platforms.

II. ART ACTIVATOR PLUG-IN

As discussed above, for ART, activation functionality will be implemented as a plug-in to a generic ART daemon. The activation plug-in is responsible for activating a target process and has no persistent data to manage. A target process is an executable process containing one or more target endpoints.

In this section the activator is defined as:

- Responsible for launching target processes
- Responsible for monitoring all child processes
- Responsible for informing the Locator about any events relating to such child processes (e.g. in particular informing the Locator if a process dies).
- Responsible for Java in-process activation or other future extensions. This is achieved via the "path name" specified for target process configuration. For example, "java.exe foo"
- NOT responsible for the management of any processes launched by previous Activators.

- NOT responsible for the monitoring and care-taking of manually activated processes.
- NOT responsible for informing specific interested parties (*ProcessMonitor*) about any process specific events.
- 5 • NOT responsible for any process check-pointing. Check-pointing is done within the Locator.
- NOT required to kill child processes if a *SIGTERM* is received.
- Not responsible for the direct management of process references (namely the immediate deletion of process references that have died).

10 Further, it is assumed that:

- Once an Activator starts up, it is able to register its endpoints with a Locator. It may need to register with multiple configured locators.
- If the Locator goes down during a process launch (or process kill), and the Activator cannot respond to the Locator, then Activator only
15 logs (audit) the fact that it could not respond to the Locator's launch request. It is assumed that the Locator will be responsible for deleting Process objects.
- Only the locator that called launch will have the object reference for the Process and so only that locator can invoke the *kill()* operation.
- 20 • The Activator reads a configuration file that defines the whereabouts of allowed Locators (i.e. the IOR of the Locators).
- There will be a configuration option to decide how long a Process in the *NOT_STARTED/DIED* state should stay in the activator. That the Activator will clean-up any old process object references whose time
25 of death is greater than the *NOT_STARTED/DIED* values.
- If an Activator dies, then its child processes do NOT necessarily die.
- The Locator sending a request to start a process that is already running is a Locator problem.
- The Activator will only accept requests to create a process from a
30 locator that the activator knows about from its own configuration. So, the activator's configuration will have a list of locator IORs to register with.

1. Object Model

(1) Object Descriptions

Referring to FIG. 5, the Activator plug-in of the present invention includes the following:

5 *IT_ActivatorPlugIn*, 101

This object represents an extension of the *main()* function in the Generic Server that is specific to the *Activator* plug-in. It is responsible for creating an *Activator*-specific transient POA from the *RootPOA* and creating and registering the *Activator* servant with it. In addition, this object registers the *Activator* with one or more *Locator(s)*.

10

IT_ActivatorImpl, 103

This is the implementation object for the *Activator* IDL interface. Only one instance of this object exists. It is responsible for creating the main *Activator* objects and handling requests for target process activation.

15

IT_ProcessTerminationMonitor, 105

This object is known as the “monitor thread” and is responsible for monitoring the processes launched by the *Activator*.

20 *IT_ProcessTerminationMonitorPosix*, 107

For POSIX, an instance of this object is created with every *IT_ProcessTerminationMonitor* creation. For POSIX, there is only one “monitor thread”.

IT_ProcessTerminationMonitorWin32, 109

25 For Win32, an instance of this object is created with every *IT_ProcessTerminationMonitor* creation. For Win32, there may be one or more “monitor thread(s)”.

IT_ProcessImpl, 111

30 This is the implementation object for the *Process* IDL interface. An object reference is passed back to the target process activation requestor so it can invoke *Process::remove()* on it. It is passed to the monitoring activity so a monitoring activity can monitor it and handle its termination. The present invention includes POSIX and Win32 versions of this object.

IT_ProcessImplPosix, 113

For POSIX, an instance of this object is created with every new *IT_ProcessImpl* creation.

IT_ProcessImplWin32, 115

- 5 For Win32, an instance of this object is created with every new *IT_ProcessImpl* creation.

IT_ProcessLauncher, 117

Given the *IT_ProcessImpl* reference, this object is known as the “launch thread” and is responsible for launching the target process.

10

IT_ProcessMonitorTable, 119

This object contains those *Process* (es) that are being monitored by one or more “monitor thread(s)”. There is only one instance of this object. Note that a *Process* will only be in either the *IT_ProcessMonitorTable* or the *IT_ProcessHeavenTable*.

15

IT_ProcessHeaven Table, 121

This object contains those *Process(es)* that have died and are no longer being monitored. There is only one instance of this object.

20 *IT_HandleTable*, 123

For Win32 only, this object contains those *Process* (es) that are being monitored by a particular “monitor thread”. The Win32 API for monitoring handles only supports 64 handles at one time.

25

(a) IT_ActivationPlugIn class

Referring to FIG. 6, the *IT_ActivationPlugIn* object represents an extension of the *main()* function in the Generic Server that is specific to the *Activator* plug-in. It is responsible for creating an *Activator*-specific transient POA from the RootPOA and

30 creating and registering the *Activator* servant with it. In addition, this object registers the *Activator* with one or more *Locator(s)*.

The *IT_ActivatorPlugIn* 101 is derived from *ITPlugIn* 101A [PLUGIN]. Via the Generic Server, this plug-in is loaded and initialized after all ORB related plug-ins are initialized and loaded. The plug-in name used for the Activator plug-in is “*Activator*”.

5 **(b) *IT_ActivatorImpl* class**

Referring to FIG. 7, the *IT_ActivatorImpl* 103 is the implementation object for the *Activator* IDL interface. Only one instance of this object exists. It is responsible for creating the main *Activator* objects and handling requests for target process activation.

10 The *IT_ActivatorImpl* 103 is derived from *POA_Activator* 103A and is the main driver for the activation plug-in. The *IT_ActivatorPlugIn* 101 creates this singleton at *orb_attach()* time before any launcher and monitor threads are created and/or started. As a result, two threads will not attempt to create this singleton. The *IT_ActivatorPlugIn* 101 deletes this singleton at *orb_detach()* time.

15 **(c) *IT_ProcessTerminationMonitor* related classes**

Referring to FIG. 8, the *IT_ProcessTerminationMonitor class* 105 object is known as the “monitor thread” and is responsible for monitoring the processes launched by the Activator.

20 The *IT_ProcessTerminationMonitor* 105 class is derived from *IT_ThreadBody* 105A. A static factory function, *smf_create_factory()*, is used for creating the platform specific instances of this class. The *IT_ActivatorImpl* 103 class is responsible for creating this object.

25 Generically speaking, the monitor thread itself is started only if there are processes to be monitored. This object uses a mutex to determine if the monitor is running. In addition, this object adds the *IT_ProcessImpl* to the *IT_ProcessMonitorTable* when a successful target process activation occurs and enters an infinite loop waiting on certain events. In the event of process death, this object removes the *IT_ProcessImpl* from the *IT_ProcessMonitorTable* and adds it to the *IT_ProcessHeavenTable*. The *IT_ActivatorImpl* deletes this object at *orb_detach()* time.

30

IT_ProcessTerminationMonitorPosix class

For POSIX, an instance of this object is created with every *IT_ProcessTerminationMonitor* 105 creation. For POSIX, there is only one “monitor thread”. This monitor thread uses the process monitor table, *IT_ProcessMonitorTable* 110, for all target processes started by the Activator and that are currently running. The `waitpid()` system call is used to wait on process death.

IT_ProcessTerminationMonitorWin32 class

For Win32, an instance of this object is created with every *IT_ProcessTerminationMonitor* 109 creation.

On Win32, there may be one or more monitor threads. Each monitor thread manages its own handle table, *IT_HandleTable* 123, consisting of its own event handles and process handles. A handle table can be configured to monitor about 62 processes.

Note that ALL monitor thread(s) share the process monitor table *IT_ProcessMonitorTable* 119 for all target processes started by the Activator and that are currently running. So as to provide three monitor threads with one having 62 processes, another having 62 processes, and another having one process. But the process monitor table includes all 125 processes. The `WaitForMultipleObjects ()` system call is used to wait on any handle event.

(d) IT_ProcessImpl related classes

Referring to FIG. 9, *IT_ProcessImpl class* 111 class is the implementation object for the *Process* IDL interface. An object reference is passed back to the target process activation requestor so it can invoke *Process::remove ()* on it. It is passed to the monitoring activity so a monitoring activity can monitor it and handle its termination.

The *IT_ProcessImpl* 111 class is derived from *POA_Process* 111A. A static factory function, *smf_create_factory ()*, is used for creating the platform specific instances of this class. The *IT_ActivatorImpl* 103 class is responsible for creating this object at *create process ()* time.

The target process activation requestor is responsible for deleting this object. Note that this object can not be deleted if it is not in a *Process: :RUNNING* or *Process : :STARTING_PROCESS* state. This object is then removed from the *IT_ProcessHeavenTable* 121. There is also an expiration timestamp on all process objects so that the Activator itself can clean up after a configured period of time.

IT_ProcessImplPosix class

For POSIX, an instance of this object is created with every new *IT_ProcessImpl* creation. At *launch()* time, this object performs a *fork()* and *exec()*. If *fork()* failed, no *pid* exists so the process object is deleted implicitly by the *IT_ProcessLauncher* and an exception is returned to the caller. Note that a *waitpid()* operation is not performed. To determine an *exec()* failure, this class creates a unidirectional pipe for communications between parent and child processes. If *exec()* fails, the child process sends a message to the parent process.

If *exec()* fails, a *pid* exists so a *waitpid()* operation is performed to avoid being marked as <defunct>. Note that the child process calls *exit()* and an exception is returned to caller. Since the child process dies, the monitor thread ignores it since it was never in a *Process:::RUNNING* state. But the *IT_ProcessLauncher* cleans up the process object implicitly. If the *fork()* and *exec()* are both successful, the process object is now monitored by the *IT_ProcessTerminationMonitorPosix*. At *kill()* time, this object performs a *kill()* system call only if the process is in a *Process:::RUNNING* state.

IT_ProcessImplWin32 class

For Win32, an instance of this object is created with every new *IT_ProcessImpl* creation. At *launch()* time, this object performs a *CreateProcess()* system call. If *CreateProcess()* failed, no *pid* exists so the process object is deleted implicitly by the *IT_ProcessLauncher* and an exception is returned to the caller. Note that a *wait* operation is not performed. If the *CreateProcess()* is successful, the process object is now monitored by the *IT_ProcessTerminationMonitorWin32*.

At *kill()* time, this object performs a *Terminateprocess()* system call only if the process is in a *Process:::RUNNING* state.

(e) *IT_ProcessLauncher class*

Referring to FIG. 10, given the *IT_ProcessImpl* reference, this object is known as the "launch thread" and is responsible for launching the target process. The *IT_ProcessLauncher* class is derived from *IT_ThreadBody*. The *IT_ActivatorImpl* class is responsible for creating this object at *CreateProcess()* times. The *IT_ProcessLauncher* invokes *IT_ProcessImpl::launch()* to start the target process. If the *wait_for_completion*

flag is specified on `create_process()`, then a “launch thread” is started by *IT_ActivatorImpl* and the requestor is notified of target process activation via the *ProcessMonitor*.

Internal Tables

5

Details of internal tables are discussed by referring to FIG. 11.

IT_ProcessMonitorTable, 119

This object contains those *Process(es)* that are being monitored by one or more “monitor thread(s)”. Note that a *Process* will only be in either the *IT_ProcessMonitorTable* or the *IT_ProcessHeavenTable*. There is only one process monitor table in the Activation plug-in. This singleton is created at ORB attach time and deleted at ORB detach time by *IT_ActivatorImpl*. This table is implemented as a hash table whose hash key is *IT_ProcessImpl::ProcessHandle*. The *IT_ProcessImpl::ProcessHandle* is guaranteed to be unique among all “live” processes whether on a POSIX or Win32 platform.

The locking model used on the *IT_ProcessMonitorTable* forces locking and unlocking the table to be performed explicitly. This is due to the fact that attributes of the process object may be modified during remove operations.

20

IT_ProcessHeaven Table, 121

This object contains those *Process(es)* that have died and are no longer being monitored. Note that a *Process* will only be in either the *IT_ProcessMonitorTable* or the *IT_ProcessHeavenTable*. There is only one process heaven table in the Activation plug-in. This singleton is created at ORB attach time and deleted at ORB detach time by *IT_ActivatorImpl*. This table is implemented as a doubly linked list.

After the death of a target process, its *Process* object is placed on the process heaven table. The *Process* object is inserted onto the front of the process heaven table. This may help performance if the assumption is that most *Activator::create_process()* requestors will request to delete the *Process* object soon after its associated target process dies. Here, the *Process* object merely sits until it is removed by its *Activator::create_process()* requestor.

30

Once a target process dies, its *pid* is not ensured to be unique. The operating system can, at any time, re-use pids. As a result, the process heaven table may have Process objects with duplicate pids. The criteria for removing a Process object from this table is the following: pid, startup timestamp, and died timestamp. All three attributes must match.

- 5 Locking is required due to the fact that the monitor thread will be adding Process objects to this table. Note that locking is performed implicitly so that callers do not have to.

IT_Handle Table, 123

- 10 For Win32 only, this object contains those *Process (es)* that are being monitored by a particular "monitor thread". The Win32 API for monitoring handles only supports 64 handles at one time. The handle table is only required for the Win32 implementation. This table is created at ORB attach time and deleted at ORB detach time when the Termination Monitor thread is created/deleted.

- 15 This table is implemented as a compact array due to the Win32 wait API which limits the number of handles it can wait for (MAXIMUM_WAIT_OBJECTS) and the fact that no fragmentation can exist in the wait API's parameter, the handle array table. The locking model used on the *IT_HandleTable* forces locking and unlocking the table to be performed explicitly. This is due to the fact that the monitor thread will be reading the
- 20 handle table while some launch threads may be adding HANDLES to this table. One *IT_HandleTable* monitors only 62 process handles. The first two slots of the *IT_Handle_Table* are occupied by two event handles. The *AddToTable* event will always reside in slot 0. The *StopMonitor* event will always reside in slot 1. The *AddToTable* and *StopMonitor* events are never removed from the handle table until the handle table is no
- 25 longer needed.

III. IMPLEMENTATION REPOSITORY

- This section describes the internal design and administrative interface to the Implementation Repository (IR). The IR is used to register and store the static data
- 30 describing server processes and OA specific information needed by the location daemon to determine the correct server process for handling incoming location requests. This document describes both the QA independent information stored in the IR as well as the POA specific information.

The Implementation Repository is an entity used by the location daemon to maintain and manage the persistent, static, data for identifying server processes and the information needed to map an incoming location request to the appropriate server process. The data in the IR identifies such things as the executable file to use for a particular server process and the host to run it on, the command line switches to pass to the executable on process startup and the environment variables that need to be set. An IR consists of an OA independent portion, which contains information about server processes and ORB instances, and an OA specific portion that is used to map an incoming location request to the proper server process. The following sections describe how this information is managed by an IR. We need to do something when an IR instance gets loaded into the Locator to validate that Ids that are unique across the entire Locator are truly unique.

1. IR Management

Referring to FIG. 12, a location daemon 201 may have more than one IR instance associated with it at any given time. The Locator 201 uses an IRRegistry 203 for managing its associated IR instances. A given IR instance may have one or more OA specific repositories with which it is associated. These repositories manage the information required by that OA's location plug-in used to map an incoming location request to the appropriate server process. The IR uses an OARespositoryRegistry to manage these OA specific repositories. Refer to Figure 1 for the following description of these interfaces.

(1) IRRegistry Interface

The IRRegistry 203 is responsible for management of IR instances associated with the Locator 201. Each IR instance has an associated name which must be unique relative to the IRRegistry 203.

```

module IT_IR
{
    interface IRRegistry
    {
        void
        create (
            in ImplementationRepository ir
        ) raises (IRAlreadyExists) ;
    }
}

```

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```

void
remove (
    in string ir_name
) raises (IRNotExist) ;

ImplementationRepository
find (
    in string ir_name
) ;

IRList
list ( ) ;
} ;

```

The **add ()** operation registers the IR instance specified by *ir*. If an IR instance is already registered with the same name attribute as the instance being added, the *IRAlreadyExists* exception is raised. The **remove ()** operation unregisters the association with the specified *ir_name* from the registry. If there is no IR instance registered with the specified *ir_name*, the *IRNotExist* exception is raised. The **find ()** operation returns a reference to the IR instance specified by *ir name*. If there is no IR instance registered with the specified *ir_name*, a NULL object reference is returned. The **list ()** operation returns a list of all IR instances registered with the registry.

(2) ImplementationRepository Interface

An ImplementationRepository 205 is responsible for management of the OA independent information needed by the location daemon to map incoming location requests to the appropriate server process. An IR instance will have one or more OA specific repositories associated with it that are responsible for maintaining that particular OA's static location information. The IR can be a persistent object reference.

```

module IT_IR
{
    interface ImplementationRepository

```

```

{
    attribute    string                name;
    attribute    boolean               enabled;
    readonly attribute string          host;
5    readonly attribute ProcessConfigRegistry
                                   process_registry;
    readonly attribute OARepositoryRegistry
                                   oa_repository_registry;

10    void
        destroy ( ) ;
    } ;
} ;

```

15 The **name** attribute is used to identify this IR instance. It is unique relative to the IRRegistry in which it is registered. The **enabled** attribute is used to enable or disable this IR instance. When an IR instance is disabled, none of its static persistent data is used by the location daemon for processing incoming requests.

20 The **host** attribute is informational only. It indicates the name of the host machine where the IR exists. It is used as an aid to administrators in knowing where various resources are allocated. The **process_registry** attribute returns a reference to the ProcessConfigRegistry instance associated with this IR instance.

25 The **oa_repository_registry** attribute returns a reference to the OARepositoryRegistry instance associated with this IR instance. The **destroy ()** operation iterates through this IR instances list of associated *OAImpementationRepository* instances calling *destroy()* on each and the persistent data associated with this instance is deleted and removed.

(3) OARepositoryRegistry

30 An OARepositoryRegistry 207 is used by the IR to manage its set of associated OA specific repositories. It is simply a map of OA repository name to abstract *OAImpementationRepository*.

Module IT_IR

```

35 {

```


The **add ()** operation adds an instance of an *OAImplementationRepository* 207 to the list of OA repositories associated with the IR instance. The repository is mapped to the provided *oa_name*. If an OA repository is already registered with the *oa_name*, the *OARepositoryAlreadyExists* exception is raised.

The **remove ()** operation disassociates the specified OA repository from this IR instance. If the specified OA repository is not registered with this IR instance the *OARepositoryNotExist* exception is raised. The **find ()** operation returns a reference to the OA repository specified by *oa_name*. If no OA repository is registered with the specified

oa_name, a NULL object reference is returned. The **list ()** operation returns a list of all the OA repository instances associated with this IR instance. The **destroy ()** operation iterates through all of the associated OA repository instances and calls their **destroy ()** operation and disassociates it from this IR instance.

5

(4) **OAImplementationRepository**

An **OAImplementationRepository** 209 interface is a simple abstract interface to an OA specific repository. It can be a persistent object.

```

10 module IT_IR
    {
        interface OAImplementationRepository
        {
            attribute string name;
15
            void
            destroy();
        };
    };
20

```

The **name** attribute is the name used for this OA repository and must be unique relative to the *OARepositoryRegistry*. The **destroy ()** operation causes the OA repository to delete and remove all of its persistent data and to destroy itself

25 **2. OA Independent Information**

The IR contains information that is not specific to any particular OA. The generic information contained in the IR is information related to server processes and location specific ORB instance identifiers (*LocOrbId*). The following sections describe this

30

(1) **Process Information**

The IR manages information about server processes through its *ProcessConfigRegistry* 221. The registry contains *ProcessConfig* 223 instances which

specify the executable file that is to be started for the server process, which host or hosts the process will run on, command line arguments and environment variables that must be provided as well as other information.

```
5  module IT_IR
   {
       enum StartUpValue
       {
           ON_DEMAND,
10      ALWAYS,
           MANUAL,
           DISABLE_START
       } ;

15      interface ProcessConfigRegistry
       {
           ProcessConfig
           add (
               in string      name,
20      in StartUpValue start_flag
           ) raises (ProcessAlreadyExists);

           void
           remove (
25      in string name
           ) raises (ProcessNotExist);

           ProcessConfig
           find (
30      in string name
           ) ;

           interface Iterator
           {
35      boolean
               next_batch (
```

```

        in unsigned long      batch_size,
        out ProcessConfigList batch
    ) ;
} ;
5      Iterator
      list ( ) ;

      void
      destroy ( ) ;
10
    } ;
} ;

```

15 The **add ()** operation creates and registers a new process configuration record with the registry. The record specifies the server process information for the target server process specified by *name*, which must be unique relative to the registry. The *start_flag* is used to specify how and when the server process should be started. If an entry already exists for the specified *name*, the *ProcessAlreadyExists* exception is raised

20 The **remove ()** operation removes and destroys the process configuration information associated with the target server process specified by *name*. All information associated with this target process is removed from the IR's persistent storage. If no entry exists for the specified *name*, the *ProcessNotExist* exception is raised.

25 The **find ()** operation returns a reference to the process configuration record specified by *name*. If no entry exists for the specified *name*, a NULL object reference is returned. The **Iterator** interface is used to iterate through a list of *ProcessConfig* entries. The **next_batch ()** operation returns the next "batch" of entries from the list. The number of entries returned in the *batch* sequence is specified by the *batch_size* parameter. If *batch_size* is zero, all the entries in the list are returned. Fewer than *batch_size* entries may be returned if the iterator has reached the end of the list. True is returned if there are more
 30 entries remaining in the list and false otherwise.

The **list ()** operation returns an iterator for the list of all the process configuration records associated with the Locator. The *destroy ()* operation removes from persistent storage all information associated with every process configuration record registered with the registry. The *ProcessConfig* interface represents the information needed

to specify a server process. The configuration information may include multiple target hosts upon which the server process may be started. Each individual host has a separate record of startup information associated with the server process to specify how that process is to be started on the particular host. The host information is held in a *HostInformation* record, described below.

```

module IT_IR
{
    struct EnvVariable
    {
        string      name;
        string      val;

    };
    typedef sequence<EnvVariable> EnvVariableList;

    typedef sequence<string> ArgumentList;

    interface ProcessConfig
    {
        attribute name;
        attribute startup flag;

        HostInformation
        add_host(
            in string      host,
            in string      activator_name,
            in unsigned long max_nbr_of_retries,
            in unsigned long retry_interval,
            in string      pathname,
            in ArgumentList arguments,
            in EnvVariableList env_variables,
            in string      user,
            in string      group
        ) raises (HostAlreadyExists);

```

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```

void
remove_host (
    in string host
) raises (HostNotExist);

Host Information
find_host (
    in string host
) ;

HostInformationList
list_hosts ( ) ;

OrbConfigList
list_orbs ( ) ;
} ;

typedef sequence<ProcessConfig> ProcessConfigList;

} ;

```

The **name** attribute specifies the unique name, relative to the ProcessConfigRegistry in which it is registered, by which this target process is referenced in other persistent records.

The **startup_flag** attribute specifies how and when this server process will be started; **ON_DEMAND** indicates that the server process will be started when a request for an object implemented by the server process arrives, **ALWAYS** is an indication that the Locator should start the server process when the daemon is started and attempt to restart it if it dies, **MANUAL** indicates that the server process will be started by an entity other than the Locator and **DISABLE_START** is an indication that the Locator should make no attempt to start that particular server process (this is a temporary state).

The **add_host ()** operation adds a new target host and its associated startup information to the list of host information for this server process instance. If an entry already exists for the specified *host*, the *HostAlreadyExists* exception is raised. The host information is described in more detail below in the description of the *HostInformation* interface. The **remove host ()** operation removes the host information for the specified *host*

and deletes all persistent storage information associated with it. If there is no entry for the specified *host*, the *HostNotExist* exception is raised. The **find_host ()** operation returns a reference to the host information for the specified *host*. If the entry does not exist a NULL object reference is returned. The **list_hosts ()** operation returns a list of the host information for all target hosts where this server process instance may be started.

The **list_orbs ()** operation returns a list of the ORB instances that are instantiated in this server process. The HostInformation 225 interface specifies the information needed to start a target server process on a specific host machine.

```

10 module IT_IR
    {
        interface HostInformation
        {
            attribute string          host ;
            attribute string          activator_name;
15         attribute unsigned long    max_nbr_of_retries;
            attribute unsigned long    retry_interval;
            attribute string          pathname;
            attribute ArgumentList     arguments;
            attribute EnvVariableList  env_variables;
20         attribute string          user;
            attribute string          group;

            void
25         set_info (
                in string          host,
                in string          activator_name,
                in unsigned long    max_nbr_of_retries,
                in unsigned long    retry_interval,
30         in string          pathname,
                in ArgumentList     arguments,
                in EnvVariableList  env_variables,
                in string          user,
                in string          group
35         ) ;
    }

```

15 } ;

20

25

30

35

“batch” update of the host information persistent data. The `get_info ()` operation provides a mechanism for doing a “batch” read of the host information persistent data.

(2) ORB Information

The IR maintains information about ORB instances that map a given ORB instance to a specific target server process. These ORB instances are referenced by a *LocOrbId* identifier that is unique relative to the Locator. To manage the ORB instance information the IR uses an `OrbConfigRegistry` 227. It should be noted that an OA repository observer can be added to the present invention.

As discussed above, an ID for an ORB instance that is unique across the entire location domain can be provided in the present invention. The ID can be treated in a hierarchical manner. In other words, a combination of the IR name and *LocOrbId* (unique relative to a given IR instance) can be used for this ID.

```

15 module IT_IR
    {
        interface OrbConfigRegistry
        {
            OrbConfig
20         add(
                in LocOrbId      id,
                in string        target_process
            ) raises (OrbAlreadyExists);

25         void

            remove (
                in LocOrbId id
            ) raises (OrbNotExist);

30         OrbConfig
            find (
                in LocOrbId id
            );
35

```

```

interface Iterator
{
    boolean
    next_batch(
5          in unsigned long batch_size,
          out OrbConfigList batch
    );
};
Iterator
10 list ();

void
destroy ();
};
15 };

```

The **add ()** operation constructs an ORB configuration record for the specified *id*, which represents an instance that will be active in the specified *target_process*. If an entry already exists for the specified *id*, the *OrbAlreadyExists* exception is raised. The **remove ()** operation removes information about the specified ORB instance from the IR. If the specified entry does not exist, the *OrbNotExist* exception is raised. The **find ()** operation returns a reference to the configuration information for the specified ORB instance. If no entry is found for the specified ORB a NULL object reference is returned. The **Iterator** interface is used to iterate through a list of *OrbConfig* entries.

The **next_batch ()** operation returns the next “batch” of entries from the list. The number of entries returned in the *batch* sequence is specified by the *batch size* parameter. If *batch_size* is zero, all the entries in the list are returned. Fewer than *batch_size* entries may be returned if the iterator has reached the end of the list. True is returned if there are more entries remaining in the list and false otherwise. The **list ()** operation returns an iterator for the list of all the ORB configuration records for this IR instance. The **destroy ()** operation removes and destroys the persistent data for all of the ORB configuration records for the IR.

An OrbConfig 229 interface represents the information maintained by the IR about a given ORB instance. Its primary responsibility is to map the ORB instance to the server process in which it will be active.

```

module IT IR
{
    typedef sequence<octet> LocOrbId;
5
    interface OrbConfig
    {
        attribute LocOrbId id;
        attribute string target_process;
10    };
    typedef sequence<OrbConfig> OrbConfigList;
};

```

The **id** attribute is an identifier for the ORB instance that is unique relative to
 15 this IR instance. The **target_process** attribute specifies the name of the target server
 process where this ORB instance will be active. It should be noted that a reference to a
 process not specified in this IR instance can be allowed.

(3) Change Notification

20 OA specific registries may need to be notified of changes in the OA
 independent information in an IR instance. To support this an Observer pattern is used.
 The following describes these Observer interfaces and how they are used.

```

module IT_IR
25 {
    interface ProcessConfigRegistry
    {
        interface Observer
        {
            void
30          add_notify (
              in ProcessConfig new_process
          );

            void
35          remove_notify (

```

```

                                in string process_name
                                );
                                };

5      typedef unsigned long ObserverId;

                                ObserverId
                                register_observer (
                                in Observer new_observer
10      );

                                void
                                unregister_observer
                                in ObserverId id
15      ) raises (
                                ObserverNotRegistered
                                );
                                // remainder of ProcessConfigRegistry interface...
                                };

20  };

```

The *ProcessConfigRegistry::Observer* interface is used by an GA specific JR to register interest in changes to the ProcessConfigRegistry.

The **add_notify ()** operation is called when a new ProcessConfig entry is added to the registry. A reference to the newly added entry is provided to the Observer in the *new_process* parameter. The **remove_notify ()** operation is called when a *ProcessConfig* entry is removed from the registry. The *process name* of the removed process is provided to the Observer. The **register_observer ()** operation is used by the OA IR to register an Observer instance with the IR. An ObserverId is returned to the caller to identify the Observer instance in subsequent calls to the registry.

The **unregister_observer ()** operation is used by the OA IR to unregister the Observer instance specified by *id* with the IR. The Observer instance will no longer be notified of *ProcessConfig* changes in the registry. If no Observer is registered with the specified *id*, an *ObserverNotRegistered* exception is raised.

Observers can be notified on individual changes to *ProcessConfig* instances.

```
module IT_IR
```

```
{
```

```
    interface OrbConfigRegistry
```

```
    {
```

```
5        interface Observer
```

```
        {
```

```
            void
```

```
            add_notify(
```

```
                in OrbConfig new orb
```

```
10            ) ;
```

```
            void
```

```
            remove_notify (
```

```
                in LocOrbId orb_id
```

```
15            );
```

```
        };
```

```
        typedef unsigned long ObserverId;
```

```
20        ObserverId
```

```
        register_observer
```

```
            in Observer new_observer
```

```
        );
```

```
25        void
```

```
        unregister_observer (
```

```
            in ObserverId id
```

```
        ) raises (
```

```
            ObserverNotRegistered
```

```
30        );
```

```
        // remainder of OrbConfigRegistry interface...
```

```
    };
```

```
};
```

35 The *OrbConfigRegistry :: Observer* interface is used by an OA specific IR to register interest in changes to the *OrbConfigRegistry*.

The **add_notify ()** operation is called when a new *OrbConfig* entry is added to the registry. A reference to the newly added entry is provided to the Observer in the *new orb* parameter. The **remove_notify ()** operation is called when a *OrbConfig* entry is removed from the registry. The *orb_id* of the removed process is provided to the Observer.

- 5 The **register observer ()** operation is used by the OA IR to register an Observer instance with the IR. An *ObserverId* is returned to the caller to identify the Observer instance in subsequent calls to the registry. The **unregister_observer ()** operation is used by the OA IR to unregister the Observer instance specified by *id* with the IR. The Observer instance will no longer be notified of *OrbConfig* changes in the registry. If no Observer is registered with the specified *id*, an *ObserverNotRegistered* exception is raised. It should be noted that
10 observers can be notified on individual changes to *OrbConfig* instances.

3. POA Specific Information

- The POA IR manages information specific to each endpoint that exists. An
15 endpoint represents either a persistent POA instance or a logical entity referred to as a namespace. The purpose of a namespace is purely administrative and thus they are not used during runtime operation of the Locator. A namespace represents a transient POA name that has been registered and thus must be unique across the location domain. By registering a namespace, an application or administrator can reserve a POA namespace hierarchy that
20 will be guaranteed to be unique across the location domain.

ART framework can provide a plug-in for the POA specific IR information, that will be loaded into the Generic Server process and work in conjunction with the Locator plug-in to provide the functionality required to support persistent POA based objects.

- 25 To manage the endpoint and namespace information the POA uses a *PoaImplementationRepository*. There may be multiple *PoaImplementationRepository* instances registered with a given *ImplementationRepository* instance.,

```

30 module IT_POA_IR
{
    typedef sequence<octet> EndpointId;

    interface PoaImplementationRepository IT_IR: :
    OaImplementationRepository

```

```

{
    NamespaceConfig
    add_namespace (
        in EndpointId id,
5        in boolean allow_dynamic_create
    ) raises (EntryAlreadyExists);

    EndpointConfig
    add_endpoint(
10        in EndpointId id,
        in boolean allow_dynamic_create,
        in LocOrbId orb_id
    ) raises (EntryAlreadyExists);

15    void
    remove (
        in EndpointId id
    ) raises (EntryNotExist);

20    NamespaceConfig
    find(
        in EndpointId id
    );

25    interface NamespaceIterator
    {
        boolean
        next_batch (
30            in unsigned long batch_size,
            out NamespaceConfigList batch
        );
    };

    interface EndpointIterator
35    {
        boolean
        next_batch (

```

```

        in unsigned long batch_size,
        out EndpointConfigList batch
    );
};

5
    readonly attribute unsigned long number_of_endpoints;

    EndpointIterator
    list_endpoints(
10        in unsigned long batch_size
    );

    readonly attribute unsigned long number of namespaces;

15
    NamespaceIterator
    list_namespaces ();

    unsigned long
    endpoints_in_process(
20        in string target_process
    );

    EndpointIterator
    list_process (
25        in string target_process
    );

    unsigned long
    endpoints_in_orb(
30        in LocOrbId id
    );

    EndpointIterator
    list_orb_id (
35        in LocOrbId id
    );
};
```


};

The **add_namespace ()** operation creates a new namespace entity in the IR to represent the POA specified by *id*. The *allow_dynamic_create* parameter indicates whether or not children of this namespace are allowed to dynamically create and register new endpoints with the Locator. If there is already an entry for either a namespace or endpoint with the specified *id*, the *EntryAlreadyExists* exception is raised.

The **add_endpoint ()** operation creates a new endpoint entry in the IR to represent a persistent POA with the specified *id*. The *allow_dynamic_create* parameter is treated the same as in the **add_namespace ()** operation. The *orb_id* indicates which ORB instance, and thus which server process, this endpoint will be associated with. If there is already an entry for either a namespace or endpoint with the specified *id*, the *EntryAlreadyExists* exception is raised.

The **remove ()** operation removes and deletes the persistent data associated with the namespace or endpoint indicated by *id*. If no entry exists for the specified *id*, the *EntryNotExist* exception is raised. The **find ()** operation returns the namespace or endpoint indicated by the provided *id*. If no entry exists for the specified *id*, a NULL object reference is returned. The **NamespaceIterator** and **EndpointIterator** interfaces are used to iterate through a list of their respective types. The **next_batch ()** operation returns the next "batch" of entries from the list. The number of entries returned in the *batch* sequence is specified by the *batch_size* parameter. If *batch_size* is zero, all the entries in the list are returned. Fewer than *batch_size* entries may be returned if the iterator has reached the end of the list. True is returned if there are more entries remaining in the list and false otherwise.

The **number_of_endpoints** attribute species the number of endpoints registered with the IR. The **list_endpoints ()** operation returns a iterator for the list of all the endpoints that are registered with the IR. The **number_of_namespaces** attribute species the number of namespaces registered with the IR. The **list_namespaces ()** operation returns an iterator for the list of all the namespaces (this includes endpoints) that are registered with the IR. The **endpoints_in_process ()** operation returns the total number of endpoints in the IR that are specified to be activated in the specified *target_process*. The **list_process ()** operation returns an iterator for all the endpoints registered with the IR that are associated with the indicated server process. NOTE: namespaces have no server process

association and thus are not included in this list. The **endpoints_in_orb ()** operation returns the total number of endpoints in the IR that are specified to be activated in the specified ORB instance. The **list_orb_id ()** operation returns an iterator for all the endpoints registered with the IR that are associated with the indicated ORB instance. NOTE:

5 namespaces have no server process association and thus are not included in this list.

(1) Namespace Information

A namespace is a logical entity that represents a transient POA instance that is used to reserve a namespace in the POA name hierarchy. It also enables the ability to

10 move endpoints from being serviced by one server process to another (see the Scenarios section below). The id for a namespace must be unique relative to the Locator.

The namespace ID can be scoped to the IR instance and use the IR name in conjunction with the namespace ID to create a unique ID relative to the locator.

```

15 module IT_POA_IR
{
    interface NamespaceConfig
    {
        attribute EndpointId id;
        attribute boolean allow_dynamic create;

        NamespaceConfig
        parent ();

        NamespaceConfigList
        children ();

    };
    typedef sequence<NamespaceConfig> NamespaceConfigList;
30 };

```

The **id** attribute identifies this namespace instance and must be unique relative to the Locator. It is a representation of the fully qualified POA name of the transient POA this namespace instance represents. The **allow_dynamic create** attribute indicates whether or not children namespaces or endpoints are allowed to dynamically

35 create and register endpoints (which represent persistent POAs).

The **parent ()** operation returns a reference to the parent namespace instance of this namespace. If this namespace does not have a parent that is registered a NULL object reference is returned. The **children ()** operation returns a list of all the registered children namespace instances of this namespace.

5

(2) Endpoint Information

An endpoint is a specialization of a namespace. It represents a persistent POA and thus may be a target endpoint for activation of persistent objects. Through its association with a particular ORB instance it specifies which server process is used to service requests for its associated objects.

10

```
module IT_POA_IR
{
    interface EndpointConfig : NamespaceConfig
    {
        attribute LocOrbId orb_id;
    };
    typedef sequence<EndpointConfig> EndpointConfigList;
};
```

20

The `orb_id` attribute specifies the ORB instance, and thus server process, where this endpoint is to be activated.

4. Scenarios

25

This section describes what happens in the IR during various scenarios of POA configuration and movement. FIG. 14 shows an example of moving a POA from one ORB instance to another in the same process and moving to a different ORB instance in another process.

30

In the first two scenarios, two persistent POAs, “B” and “C”, that are children of a transient POA “A” and are in the same ORB instance. Each POA has an endpoint entry in the IR’s endpoint table that maps the endpoint ID to the LocOrbId for the ORB instance and there is an entry in the ORB table that maps the LocOrbId to a process. To move POA “C” from one ORB instance to another, a new entry must be created in the ORB table that maps the new LocOrbId to “Process 1” and the entry for POA

"C" in the endpoint table must be changed to map to the new LocOrbId. Moving POA "C" to another process is very similar to moving it to another ORB instance. An new entry is created in the ORB table for a new LocOrbId which is mapped to "Process 2", and the entry in the endpoint table is modified to map to the new LocOrbId. From this example it should be evident that moving all the endpoints associated with an ORB instance from one process to another is as simple as modifying the entry for their associated LocOrbId in the ORB table to point to the new process.

Scenario 3 of FIG. 14 shows what happens when a POA with a persistent parent POA is moved to another process. In this case the moved POA "C" will have a transient parent POA "A", since the objects associated with POA "A" must be active in "Orb I" of "Process 1". It is not necessary in this situation to have a namespace entry for POA "A", even though it is logically a namespace, since an endpoint entry already exists for POA "A", thus reserving the POA "A" namespace.

IV. Generic Server

ART is designed for pluggability. There may be many types of plug-ins. This section addresses application-level plug-ins. An application-level plug-in is defined as a plug-in that is tied to the lifetime of a target process, whereas an ORB-level plug-in is tied to the lifetime of an ORB instance. An application-level plug-in can provide a locatability service, a process activation service, a naming service, etc. The application-level plug-ins provide the flexibility for a wide range of configuration and installation scenarios. How to configure application-level plug-ins will be left up to an installation.

The generic ART server is a container for an application-level plug-in. The generic ART server acts as the *main()* for an application-level plug-in. One or more generic ART servers may be configured within one location domain.

Note that services provided in a location domain are *not* required to be implemented as plug-ins and are not required to use the generic ART server. But some of the benefits are:

- Allows component-izing services within ART.
- Allows fewer executables running on a system. It should be noted that more than one application-level plugin in a generic server can be provided. It should also be noted that the generic ART server is closely tied to ORB Management and Plug-In Management.

Furthermore, the generic ART server is itself a target process in the ART framework, and all plug-ins are loaded and initialized on a per ORB basis.

FIG. 15 illustrates two horizontal views of the generic server model. One from location domain to administrative domain, and then to ORB instance. The second
 5 from host to active generic server to ORB instance. One or more ORB instances may be initialized in one administrative domain spanning one or more hosts and generic servers. In addition, one or more plug-ins may be loaded and initialized in one ORB instance.

As noted above, the generic server acts as the *main ()* for a single application-level plug-in. The application-level plug-in name is specified as an argument to
 10 the generic server. The generic server performs the following:

- Initializes the ORB via *CORBA::ORB_init()* with ORB-level plug-ins.
- Activates the root POA manager so that Application-level plug-ins can create their own POAs and POA managers from the root POA.
- Obtains a reference to the *ConfigManager*.
- 15 • Obtains a reference to the *PlugInManager*.
- Loads and registers the Application-level plug-in(s) specified at the command line via *PlugInManager::getplugin()*.
- Initializes the Application-level plug-in(s) specified at the command line via *ITAppPlugIn::app_init()*.

20 To shutdown the generic server, send a CTRL-C signal. The generic server's signal handler performs the following:

- Shutdown the application-level plug-in via *ITAppPlugIn::app_shutdown()*.
- Shutdown the ORB and its ORB-level plug-ins via *CORBA::ORB::shutdown()*.

It should be noted that a plug-in controlled POA with its own POA manager
 25 is created. In this way, all objects created will be managed by its own POA.

The generic server can support one or more ORB instances, and can support one or more application-level plug-ins.

The following features can also be proved in the generic server:

- Command line switches as well as an environment setting for the generic ART
 30 daemon.
- Relationship to an In-process Activator; Launching servers in separate processes.

- An app-level plug-in can be designed to either be stand-alone or dependent on other plug-ins. An app-level plug-in loaded into a generic server can locate objects that might live in other app-level plug-ins.
- The generic server can be dynamically signaled to load or unload a certain plug-in.
- 5 A list of currently-loaded plug-ins can also be provided.
- Each app-level plug-in can use its own POA.
- The plug-ins are required to be dynamically loaded as well as statically linked to a generic ART server. The static linking requirement not defeat the purpose of a generic server. Alternatively, either static linking or dynamic loading to build the
- 10 generic server can be provided.

V. General Plug-In

This section describes plug-ins in general and how they are managed by the ART core. It is expected that there will be a number of different types of plug-ins. All

15 types, however, are derived from the *ITPlugIn::PlugIn* interface and share a common functionality.

All plug-ins are preferably managed by a single *ITPlugIn::PlugInManager* object within the ART core.

1. An Overview of Plug-Ins

20 As noted above, there is a base plug-in interface, *ITPlugIn: PlugIn*. This interface contains common plug-in information, like the plug-in's name and version information. Plug-ins can be loaded and unloaded, registered and unregistered.

25 Loading a plug-in means loading the module in which the plug-in resides. A plug-in can be loaded into an application in a number of ways:

- Linked directly to the application, either as object code or in a static library.
- Linked with the application as a dynamic library and loaded by the system on start-up.
- Loaded explicitly by the ORB or application when required.
- 30 • Loaded implicitly by being contained in the same module as an explicitly loaded plug-in.

After a plug-in is loaded, it is registered with the *PlugInManager* to make it available to the application. The plug-in developer call *PlugInManager::register_plugin ()* after the plug-in has been instantiated. This can be done from the plug-in's constructor.

The plug-in model is as follows:

5 The plug-in implementation code resides in a library. When the library is loaded, the plug-in is instantiated and its constructor registers it with the *PlugInManager*.

Plug-ins can be unloaded. Note that only dynamically loaded plug-ins can be unloaded. Unloading a plug-in causes it to be unregistered also.

10 2. Plug-in Initialization

After a plug-in is loaded and registered it is available for use by the application. However, different types of plug-in are used in different ways and at different times. To use a plug-in, the application first initializes it. When the application is finished using a plug-in, it informs it of this.

15 It is left to derived plug-in types to define their initialization and shutdown behavior. For example, an ORB level plug-in needs to be initialized and shutdown for each ORB instance, whereas an Application level plug-in is usually only and shutdown once.

Further, plug-ins are identified by their names. A plug-in's name is simply a string. As such, there is the possibility for name clashes. A naming convention can be
20 provided in order to minimize name clashes.

In addition, it is somewhat specific to a derived plug-in type as to how arguments are passed to a plug-in in it's initialization phase. Therefore, an argument convention can be provided for argument naming, behavior of modifying the argument list, etc.

25 Plug-ins are expected to use the core configuration mechanism. A plug-in should retrieve its configuration information from a config scope of the same name as the plug-in. For example, the IIOP plug-in would expect to find values for its configuration variables in the "IIOP" configuration scope.

Moreover, a plug-in may require other plug-ins to be available before it can
30 successfully be used. It can list these dependencies in the reserved configuration variable *PrerequisitePlugIns*, which may be present in any plug-in's configuration scope. The value for this variable is simply a list of plug-ins that must be loaded before this one. The *PlugInManager* will attempt to load each plug-in in this list. If one of these plug-ins cannot

be loaded then the dependent plug-in will not be loaded, and a *ITPlugIn : : LoadFailed* exception is thrown. This also raises the possibility of circular dependencies, e.g., if plug-in A depends on plug-in B, and plug-in B depends on plug-in A. The *PlugInManager* will detect such circular dependencies and raise a *ITPlugIn : : CircularDependency* exception.

5 An implementation of the *ITPlugIn : : PlugInManager* interface returns an object reference to the *PlugInManager* implementation object. This object reference is released after use.

The *PlugInManager* is used to load, unload, register and unregister plug-ins. But its main function is to obtain a reference to a plug-in.

10 When *PlugInManager : : get_plugin ()* is called, the *PlugInManager* first checks its list of registered plug-ins. If a plug-in of that name is registered, an object reference to the plug-in is returned. If that plug-in has not already been registered, the *PlugInManager* uses the *ConfigManager* to get the value for the "LibraryName" configuration variable in the plug-in's configuration scope, and attempts to load it. It is
15 assumed that the plug-in is registered with the *PlugInManager* when it is loaded. If the plug-in is successfully loaded but not registered, *PlugInManager : : get_plugin ()* throws a *NotRegistered* exception. If the plug-in is successfully loaded and registered, a reference to it is returned.

20 The caller of *PlugInManager : : get_plugin ()* can then narrow the returned reference to the appropriate type.

3. Interfaces

The following interfaces are locality constrained.

```
25  //      IDL
    module ITPlugIn
    {
        typedef unsigned short PlugInType;
        const PlugInType ORB_PLUGIN = 0;
30      const PlugInType APP_PLUGIN = 1;

        interface PlugIn
        {
35          readonly attribute PlugInType type;
```



```
        readonly attribute string name;

        // More general plugin information.
    } ;
5  interface PlugInManager
    {
        exception LoadFailed
        {
            string reason;
10
        } ;

        exception CircularDependency { } ;
        exception UnloadFailed { } ;
15        exception NotRegistered { } ;
        exception AlreadyRegistered { } ;
        exception UnregisterFailed { } ;

        PlugIn
20        get_plugin (
            in string plugin_name
            in ITCfg : : ConfigManager config
        ) raises (LoadFailed, CircularDependency, NotRegistered) ;

25        void
        unload_plugin (
            in string plugin_name
        ) raises (UnLoadFailed);

30        void
        register_plugin (
            in PlugIn plugin
        ) raises (AlreadyRegistered) ;

35        void
        unregister_plugin (
            in string plugin_name
        ) raises (UnregisterFailed) ;

    } ;
```

} ; The plug-ins can be configured to be dynamically upgraded, or unloaded without shutting down the process.

VI. ORB Plug-In

5 This section describes ORB level plug-ins. An ORB level plug-in is one used by the ORB core. It is associated with an ORB instance.

Each ORB instance in a process is a member of an administrative domain. One of the functions of this domain is to specify what ORB level plug-ins are required for any ORBs that are members of it. This is done through a configuration variable called
10 RequiredORBPlugIns.

When *CORBA::ORB_init()* is called to initialize an ORB instance, each of the ORB's required plug-ins is also initialized. The ORB uses the *PlugInManager* to find and load any required plug-ins. It then calls each plug-in's *ORBPlugIn::ORB_init()* method, passing in itself. If plug-in A is dependent on plug-in B, then plug-in B (and any of
15 it's dependencies) is initialized before plug-in A.

If a required plug-in, or any plug-in it's dependent on, cannot be successfully loaded and initialized, *CORBA::ORB_init()* raises a *CORBA::INITIALIZE* exception.

It should be noted that when the *CORBA::INITIALIZE* exception is raised, corresponding errors can be reported.

20 During its *ORBPlugIn::ORB_init()* call, the plug-in should create a *PerORBState* object. This object can hold any ORB-specific objects that the plug-in creates. This *PerORBState* object is returned to the ORB, which stores it.

The ORB level plug-ins can be configured purely through the ORB, i.e., its *ConfigManager*. This would allow the plug-in readonly access to the ORB's arguments if
25 necessary, but discourages passing argument's to plug-ins directly.

It should be noted that not all plug-ins are required to be loaded before the ORB can operate successfully. Some plug-ins may be available to the ORB, but rarely used. In such cases, these plug-ins are not loaded until they are needed. For example, a security plug-in may not be needed until an encrypted request arrives.

30 One way a plug-in may be loaded on demand is by calling *PlugInManager::get_plugin()*, passing in the name of the plug-in to be loaded.

Another way is through *CORBA::ORB::resolve_initial_references()* (in *ObjectId* identifier). This method provides a way for applications to obtain references to

various objects and services, identified by an *ObjectId*. If the required object or service lives in a plug-in, then `CORBA::ORB::resolve_initial_references()` will cause the plug-in to be loaded and initialized. During the plug-in's initialization phase, it will register the reference, which is then returned.

This works as follows: First, the *ConfigManager* is used to get a value for the “PlugInName” variable in the *<ObjectId>* scope. For example, a configuration file might contain the line, *POA.PlugInName = “POApi”*. This plug-in is then loaded via the *PlugInManager* and initialized. As a consequence, a POA initial reference is registered, and returned to the caller.

Furthermore, when the ORB^{*} is shutting down it first calls *PerORBState* : : *ORB_shutdown_phase1* () and then *PerORBState* : : *ORB_shutdown_phase2* () on each of its stored *PerORBState* objects. The ORB plug-in can then clean up any objects held by the *PerORBState* object.

The following is a exemplary interface:

```
// IDL
module ART_Plugin
{
    interface PerORBState
    {
        void
        ORB_shutdown_phase1 ( ) ;

        void
        ORB_shutdown_phase2 ( ) ;
    } ;

    interface ORBPlugin
        : ITPlugin : : Plugin
    {
        PerORBState
        ORB_init(
            in ORB orb
        )
    } ;
} ;
```

VII. Event Handling

This section addresses event handling within ART. It presents a design for an event handler in the ART core. Event handling includes providing an interface between the ORB and the transports, allowing the main thread to be allocated to do ORB work when necessary, demultiplexing of multiple transports and the handling of foreign event loops.

The following is a discussion of the environment in which the event handling of the present inventions operates. The ORB Portability Joint Submission adds several thread-related operations to the ORB interface. These operations are included to support single-threaded ORBs as well as multi-threaded ORBs that run multi-thread-unaware code.

A single-threaded application can make no assumptions about the underlying ORB. To remain thread-safe, it specifies the Single Thread POA policy. To allow the ORB to perform work, it must either call *run()* or *perform_work()*. By using the Single Thread POA policy, the application is guaranteed that requests will be processed sequentially, and any upcalls will be made in a thread-safe manner in the main thread.

A multi-threaded application that specifies the ORB Controlled POA policy does not need to call any of the above operations. However, calling **run()** may be useful to prevent the application from exiting until the ORB is shut down. A multi-threaded server may use the Single Thread POA policy, or a mixture of Single Thread and ORB Controlled policies. In this case it calls *run()* or *perform_work()* to allow work to be done by the POA with the Single Thread policy.

The new thread-related operations on the ORB are assumed to behave as follows:

boolean

ORB::work_pending()

This operation returns true if there is work to be done *and* it needs to be done in the main thread. In a single-threaded ORB all work will be done in the main thread. A multi-threaded ORB will do work involving a POA with the Single Thread policy in the main thread, while ORB Controlled POAs may do work in any thread.

void

ORB::perform_work()

If called by the main thread, does some work. Otherwise does nothing.

void

ORB::run()

If called by the main thread, enables the ORB to do work using the main thread, until the ORB is shut down. Otherwise waits until the ORB shuts down. This does not mean the ORB can do not work, just that the thread the work is done in is undefined. A consequence of this is that no work associated with POAs that have specified the Single Threaded policy will be done if *run()* is called in a thread other than the main thread.

void

ORB::shutdown(boolean wait_for_completion)

Instructs the ORB to shut down. It optionally blocks until all ORB processing has completed. This includes request processing, object deactivation and other operations associated with object adapters. After this operation has been called, the ORB is no longer useable until *ORB_init()* is called again.

Details

The basic design of the event handler is that it contains a work queue and a Reactor described below. The interface to the event handler will be the same for both single- and multi-threaded implementations of ART, but the implementation of the event handler itself may be different.

Multi-threaded case

The work queue in the multi-threaded case is there to provide a means for work to be done in the main thread. Work relating to a POA that has the Single Thread policy can be done in the main thread.

When a request needs to be processed, the POA provided interceptor can determine which policies are in use. If the POA has the ORB Controlled thread policy, then the request gets dispatched immediately. Alternatively, if the Single Thread policy has been specified, the request gets put on the work queue to be processed by the main thread later.

The main thread can do work on the queue by calling *ORB::perform_work()* or *ORB::run()*. The implementation of these functions is to check if there is any work pending on the queue and process it if so. If these calls are made from any other thread, no work will be done.

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```

        handle_events(
            in long timeout
        );
        void
5      shutdown();
        void
        register_handle(
            in ART_EventHandler h,
            in EventType t
10      );
        void
        remove_handler(
            in ART_EventHandler h,
            in EventType t
15      );
        void
        suspend_handler(
            in ART_EventHandler h
        );
20      void
        resume_handler(
            in ART_EventHandler h
        );
25  };

```

void

ART_Reactor::handle_events(long timeout)

Starts the Reactor processing events, for *timeout* milliseconds. A *timeout* of IT_INFINITE_TIMEOUT means this method will not return until *shutdown()* has been

30 called.

void

ART_Reactor::shutdown()

Stops the Reactor. The Reactor will finish any operation it is currently in,

35 then it will exit it's *select()* loop, causing *handle_events()* to return. It removes all handles from itself, so each registered event handler will get a *handle_close()* callback.

void

ART_Reactor::register_handler(ART_EventHandler h, EventType t)

Register the handler *h*, to be called back when an event of type *t* occurs.

Event types can be READ_EVENT, EXCEPT_EVENT, WRITE_EVENT or some

5 combination of all three, by bitwise 'or' ing them together. The handler gets called back on *handler_input()* when a READ_EVENT occurs, *handle_exception()* on an EXCEPT_EVENT, etc.

void

10 **ART_Reactor::remove_handler(ART_EventHandler h, EventType t)**

The handler *h* will no longer be called for events of type *t*. If this means that the handler is not interested in any events, the handler is removed from the Reactor.

void

15 **ART_Reactor::suspend_handler(ART_EventHandler h)**

The handler does not get called when any events occur. The handler is not actually removed however, and can be reactivated by calling *resume_handler()*.

void

20 **ART_Reactor::resume_handler(ART_EventHandler h)**

The handler that was previously suspended by *suspend_handler()* is reactivated.

25 The following is an EventHandler interface:

native ART_Handle; // unsigned int - a file descriptor

interface ART_EventHandler

{

ART_Handle

30 get_handle();

void

set_handle(

in ART_Handle h

);


```

        boolean
        handle_input();
        boolean
        handle_output();
5       boolean
        handle_exception();
        boolean
        handle_timeout();
        boolean
10      handle_close();
    };

```

ART_Handle

ART_EventHandler::get_handle()

Returns the handle associated with this event handler.

void

ART_EventHandler::set_handle(ART_Handle h)

Associates the handle *h* with this event handler.

boolean

ART_EventHandler::handle_input()

Called when the file descriptor associated with the handler is ready for reading. This method should not block.

This method will also get called if the associated socket is closed by the client side. In this case there may be no data to be read from the socket. If this happens, the handler should be removed from the Reactor.

boolean

ART_EventHandler::handle_exception()

Called when the file descriptor associated with the handler is ready for writing. This method should not block.

boolean

ART_EventHandler::handle_output()

Called when the file descriptor associated with the handler is ready for writing. This method should not block.

boolean

ART_EventHandler::handle_timeout()

Called when a timeout has occurred.

boolean

ART_EventHandler::handle_close()

Called when the file descriptor associated with the handler has been removed from the Reactor. The file descriptor can now be safely closed.

The purpose of the work queue is to provide a way for the main thread to perform work associated with ORBs that specify the Single Thread policy, in a multi-threaded environment. Such work can be done in the *main* thread.

The work queue is populated by transports or interceptors. This work is processed when the main thread calls *ORB::perform_work()* or *ORB::run()*. Note that since *ORB::perform_work()* only does work associated with that ORB, there will be a work queue per ORB. However, see the busy-loop issue below for a discussion on why it may be useful to allow the ORB to specify that its work can be done when *ORB::perform_work()* is called for *any* ORB. This could be achieved by having a process wide work queue that ORBs can share.

The following is a WorkQueue interface:

interface ART_Work

```
{
    void
    do_work();
};
```

interface ART_WorkQueue

```
{
    boolean
    is_empty();
    boolean
```

```

    post_job(
        in ART_Work job,
        in short priority
    );
5   void
    process_jobs(
        in long timeout
    );
    boolean
10   process_one_job(
        in long timeout
    );

    void
15   stop_processing(
        in boolean process_remaining
    );
    void
    flush();
20 };

    void
    ART_Work::do_work()
        Called when job gets processed.

    boolean
25   ART_WorkQueue::is_empty()
        Returns true if the queue is empty.

    boolean
30   ART_WorkQueue::post_job(ART_Work job, short priority, long
        timeout)
        Posts a job to the queue at the specified priority. The priority is represented
        by a value between 0 and 31. If the timeout is 0, this method fails to post and returns
        immediately if the queue is full. A timeout of IT_INFINITE_TIMEOUT means wait
35   indefinitely to post the job.

```

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void

ART_WorkQueue::process_jobs(long timeout)

Process the queue. If *timeout* is *IT_INFINITE_TIMEOUT*, this method will not return until *stop_processing()* is called. If *timeout* is 0, it returns immediately if there is no work to be done.

boolean

ART_WorkQueue::process_one_job(long timeout)

Process the first job on the queue. If *timeout* is *IT_INFINITE_TIMEOUT*, this method will not return until one job has been processed. If *timeout* is 0, it returns immediately if there is not work to be done. Returns true if work was done, false if it timed out.

void

ART_WorkQueue::stop_processing(boolean process_remaining_

Stops the queue if it is processing jobs or waiting for jobs to *process_remaining* is true, then all jobs currently in the queue are processed. If *process_remaining* is false then all jobs in the queue are discarded. Any job that is being currently processed is completed.

void

ART_WorkQueue::flush()

Empties the queue.

To hook into the Reactor, a transport plug in performs do the following:

First, create one or more implementations of the *ART_EventHandler* interface. For example, a typical transport plug-in would have one implementation for accepting new connections on a listening socket, and another for reading data from the connected sockets. Each instance of these *EventHandler* implementations is associated with a handle, or file descriptor, and gets called back whenever anything happens on that handle. For example if the file descriptor is ready for reading, the *EventHandler::handle_input()* method gets called.

In the plug-in's *orb-initialise()* code, obtain a reference to the Reactor using *ORB::resolve_initial_references("IT_Reactor")*. The *EventHandlers* can then be registered with the Reactor, using the *register_handler()* method. The *EventType* parameter indicates

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what type of event the EventHandler is interested in being notified of by the Reactor. This can be *ART_Reactor::READ_EVENT*, *ART_Reactor::EXCEPT_EVENT*, *ART_Reactor::WRITE_EVENT*, or some combination of the three.

EventHandlers can indicate that they are no longer interested in a particular type of event by calling the *remove_handler()* method. If the EventHandler is not interested in any type of event, it's removed from the Reactor. EventHandlers can be temporarily switched off and on by the *suspend_handler()* and *resume_handler()* methods.

When an EventHandler is removed from the Reactor, it gets called back on the *handle_close()* method. The associated file descriptor can then be safely closed. The file descriptor should not be closed before this method is called, because the Reactor could still be *select()* ing on it.

It should be noted that the Reactor can be limited to file descriptor based transports in single-threaded mode.

In both single- and multi-threaded models, use of *select()* in the Reactor can limit the number of file descriptors of *MAX_FD_SETSIZE - 1*, which is typically 1023 on Unix and 63 on Win32. This limit could be overcome on Win32 by having multiple Reactors in different threads, but not on Unix due to the fact that *fd_sets* are bitsets.

The busy-loop problem can be described as follows: Consider a process containing multiple ORBs, and each ORB needs to use the main thread to perform some work. Since there is a different work queue for each ORB the application must check each one in turn to see if there is any work to be done. If there is not work to be done for any ORB, the process is still using processor time by constantly checking.

One solution to this problem is to allow an ORB to state that it wants it's work to be done whenever *ORB::run()* is called - even if it is called by another ORB. This is like an extra ORB policy. If this policy is not set, then no work would be done for that ORB until *run()* or *perform_work()* is called on it. This could be implemented by having the ORB use a process-wide work queue, instead of a per-ORB one.

The present invention allows an application to register its file descriptors with it, taking control over the applications fd based events. It is also possible for an applications to control events of the present invention by getting the file descriptors. There are two cases.

In the first case, the application can register it's file descriptors with ART. This is simply a matter of registering the foreign file descriptors with the Reactor. The

application would derive from *ART_EventHandler*, register an instance of this class with the Reactor for each of it's file descriptors, and get called back whenever one of it's file descriptors fire.

5 In the second case, the ART can register it's file descriptors with an application. In turn, there are two ways to achieve this. One way is to disable the Reactor in the ART core, and instead just hand over any file descriptors registered with it to the application. The application will then call the Reactor back whenever one of the file descriptors file, and in turn the appropriate event handler is called. The second method is to leave the Reactor running as normal within the core. When the application ask for file
10 descriptors, one end of a pipe can be returned. Then, whenever an event gets posted to the work queue, one byte can be written onto the pipe. The application will notice that an event is ready and call back. Then the event can be processed on the queue. Note that this second method only works in a multi-threaded environment, and the application only gets notification when work needs to be done in the main thread.

15 Although the preferred embodiments of the invention have been described in the foregoing description, it will be understood that the present invention is not limited to the specific embodiments described above.

THE CLAIMS

What is claimed is:

1. A computer implemented method of activating a process, comprising:
5 generating one or more first plug-ins each configured to activate a target process;
dynamically registering the first plug-ins with a second plug-in; and
permanently storing information relating to each registered first plug-in.
2. The method of claim 1 further comprising:
10 storing a flag for each registered first plug-in;
perpetually activating the corresponding target process if the flag is set to a first
state; and
activating the corresponding target process upon a request if the flag is set to a
second state.
- 15 3. The method of claim 2 further comprising:
generating an exception to indicate that a target process is inactive when its flag is
not set to the first state or the second state.
- 20 4. The method of claim 1 further comprising:
providing a unique identifier for each target process; and
sending and receiving a message between the first and second plug-ins using the
identifiers.
- 25 5. The method of claim 4 wherein the message includes information relating to a state
change of the target processes, and wherein the state includes an activated state and a
deactivated state.
- 30 6. A server computer in a client-server computer system, comprising:
one or more first plug-ins each configured to activate a target process; and
a second plug-in configured to dynamically register the first plug-ins and to
permanently store information relating to the registered first plug-ins.

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7. The server of claim 6 wherein the second plug-in comprises:

a memory configured to store a flag for each registered first plug-in, wherein the second plug-in is further configured to perpetually activate target processes having their flags set at a first state and to activate target processes, upon receiving a request, having their flags set at a second state.

8. The server of claim 7 wherein the second plug-in is further configured to generate an exception to indicate that the target process is inactive when the flag is not set to the first state or the second state.

9. The server of claim 6 further comprising:

a first computer program object configured to provide a unique identifier for each target process and configured to send a message using the identifiers.

10. The server of claim 9 wherein the message includes information relating a state change of the target processes, and wherein the state includes an activated state and a deactivated state.

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(54) Title: METHOD AND SYSTEM FOR DYNAMIC CONFIGURATION OF ACTIVATORS IN A CLIENT-SERVER ENVIRONMENT			
(57) Abstract <p>A computer implemented method of activating a process. The method includes the steps of generating one or more first plug-ins each configured to activate a target process, dynamically registering the first plug-ins with a second plug-in, and permanently storing information relating to each registered first plug-in. A server computer in a client-server computer system that includes one or more first plug-ins each configured to activate a target process, and a second plug-in configured to dynamically register the first plug-ins and to permanently store information relating to the registered first plug-ins.</p>			

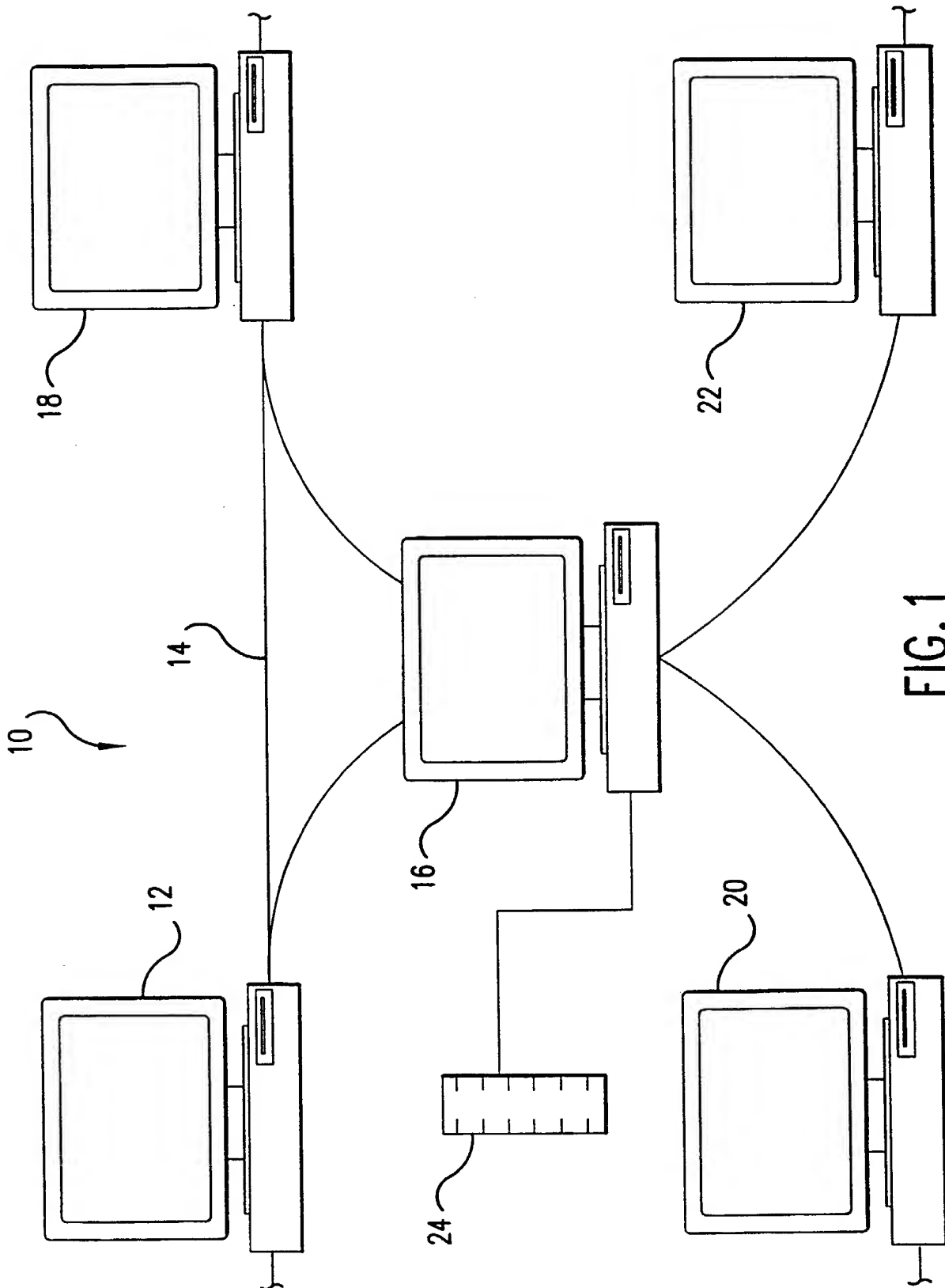


FIG. 1

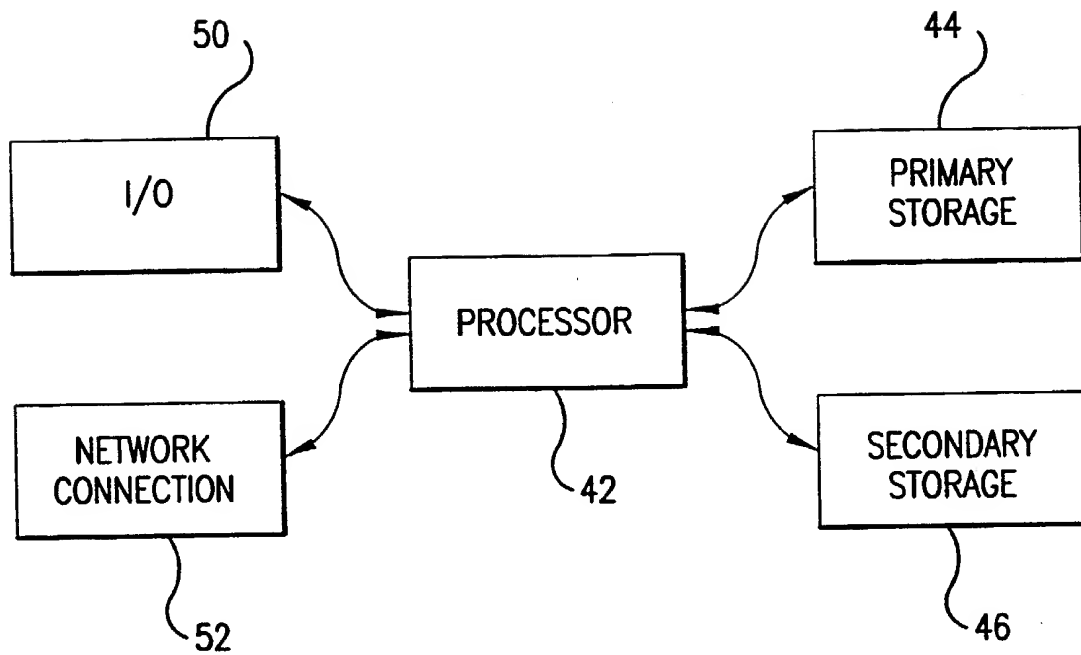


FIG. 2

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ACTIVATOR AND LOCATOR COMMUNICATIONS

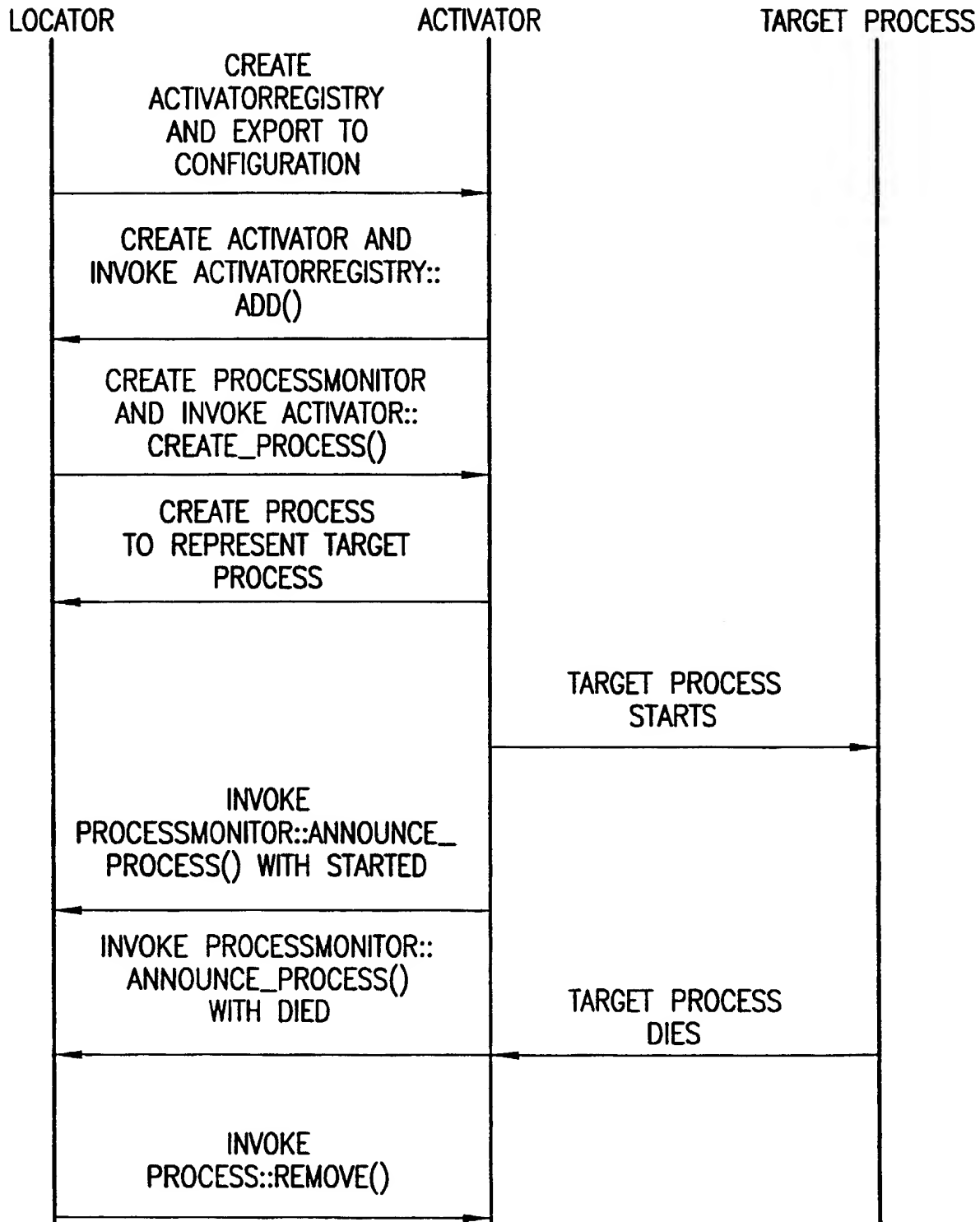


FIG. 3

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ACTIVATION AND LOCATION CLASS ASSOCIATIONS

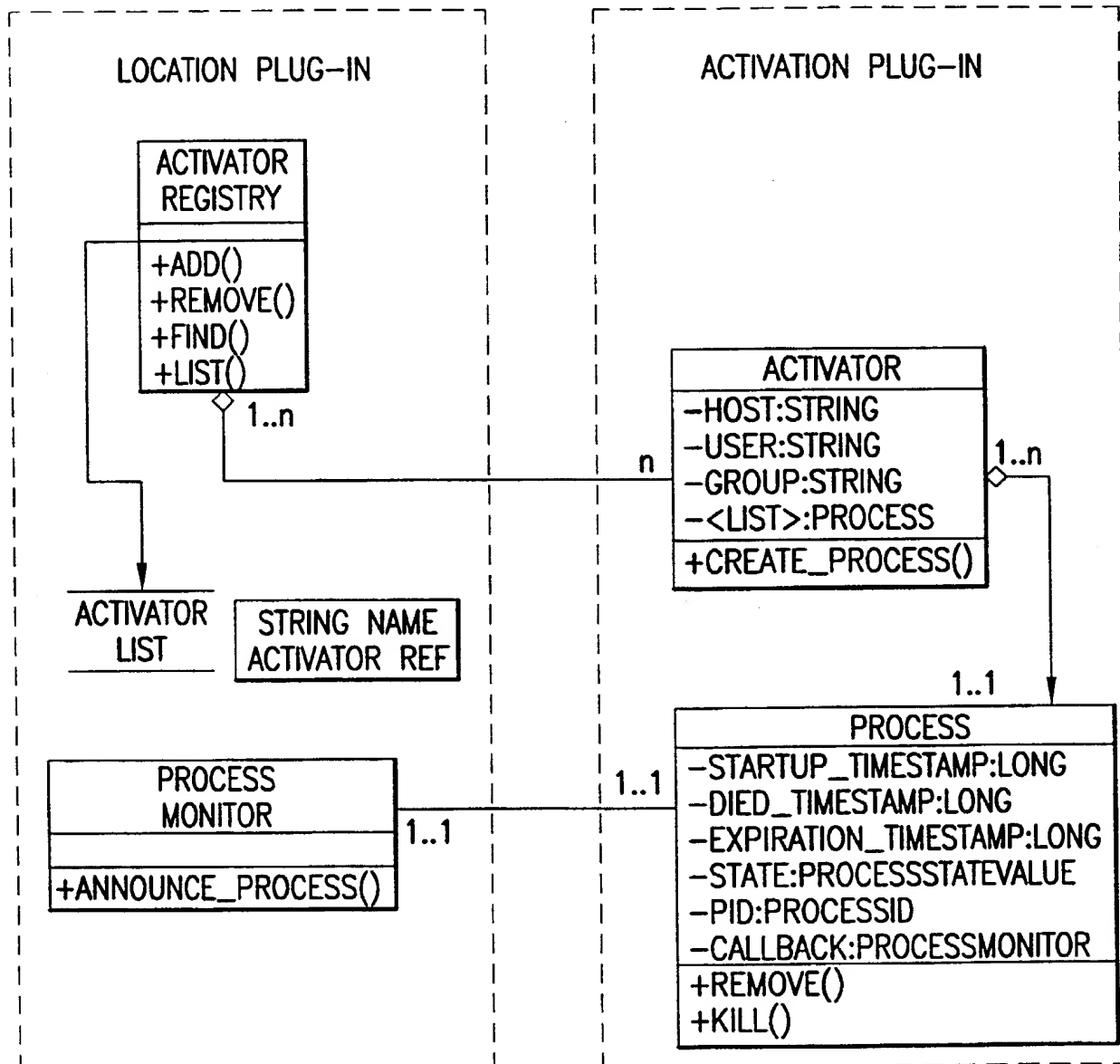


FIG. 4

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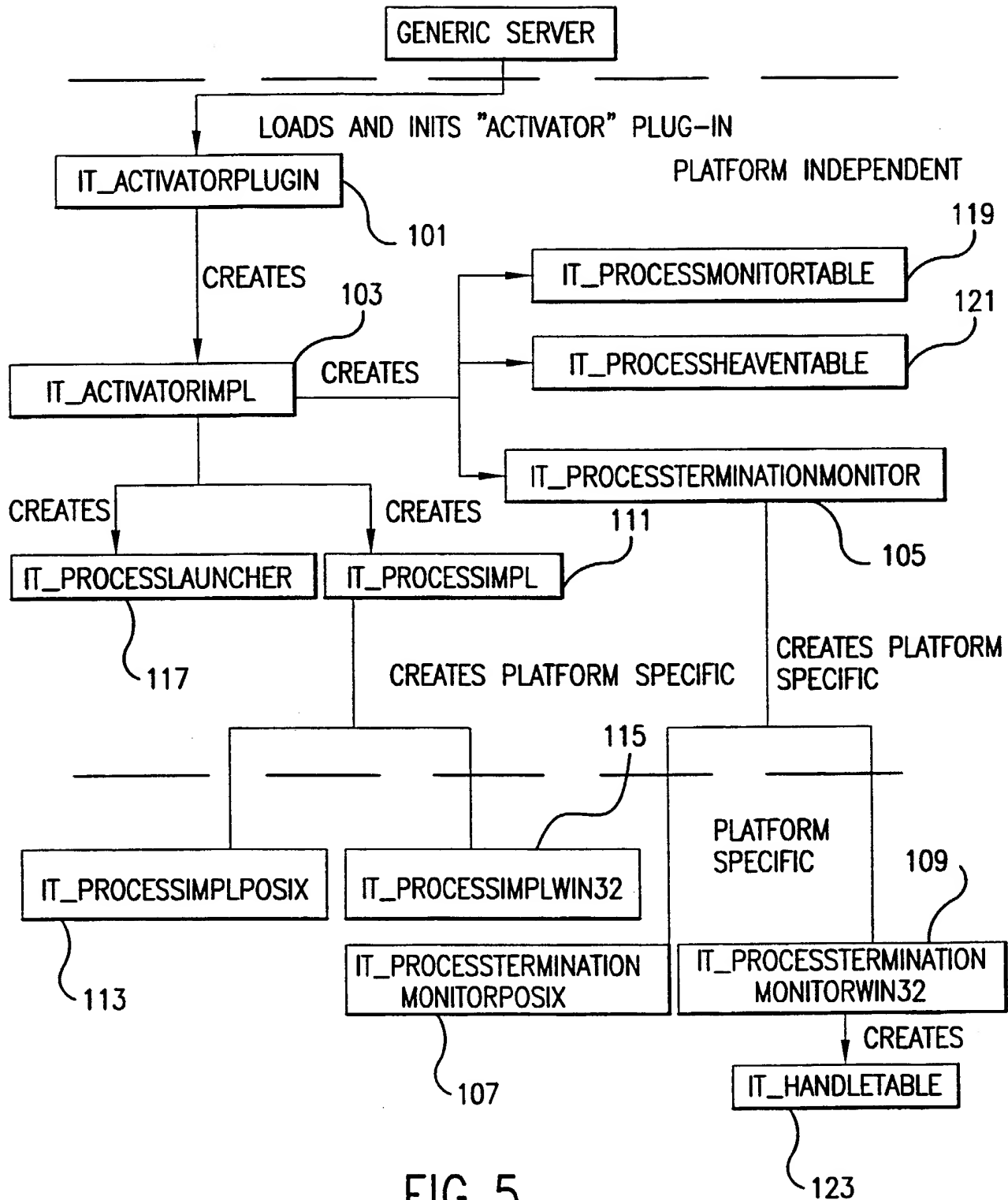


FIG. 5

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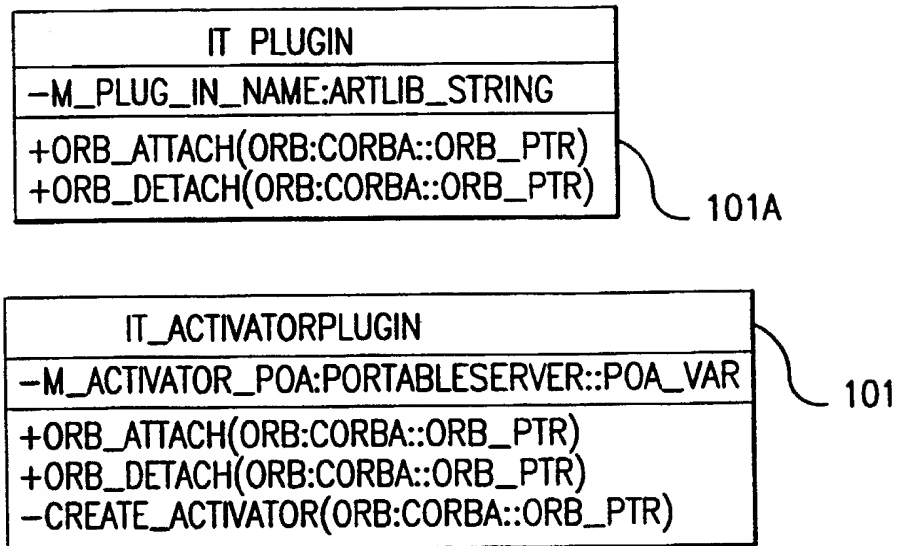


FIG. 6

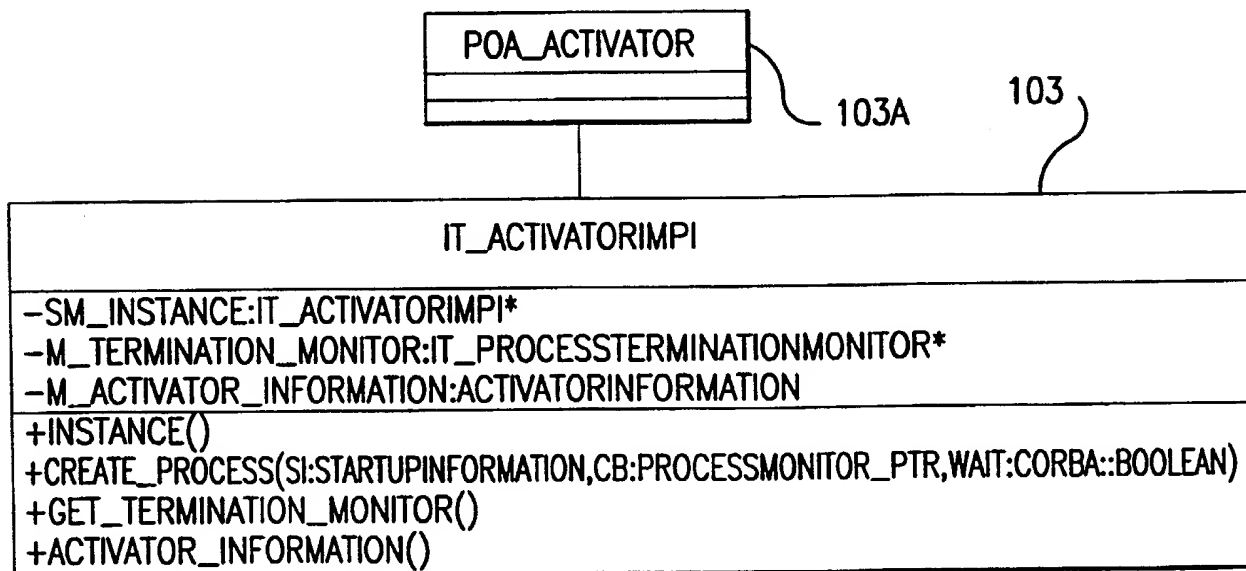


FIG. 7

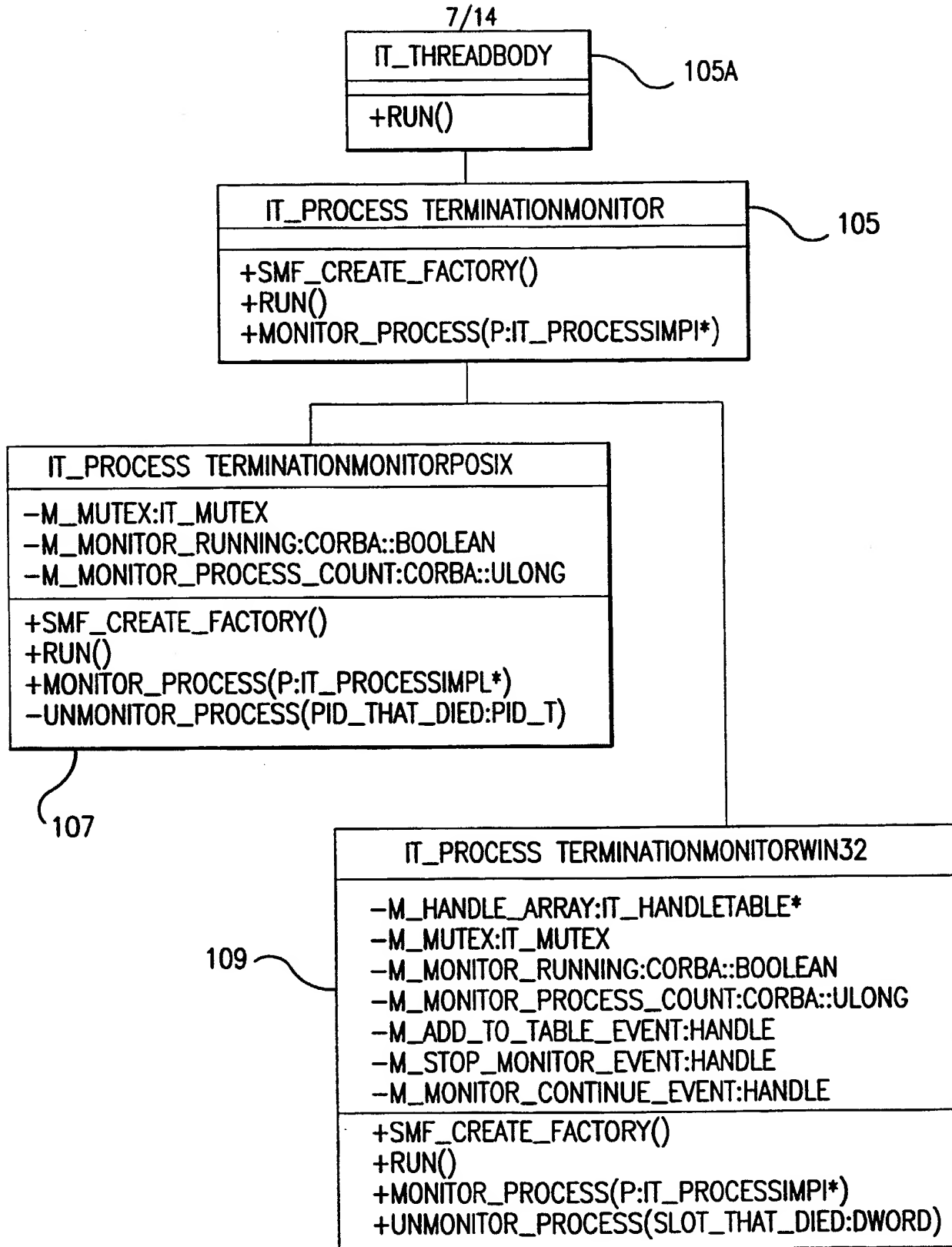


FIG. 8

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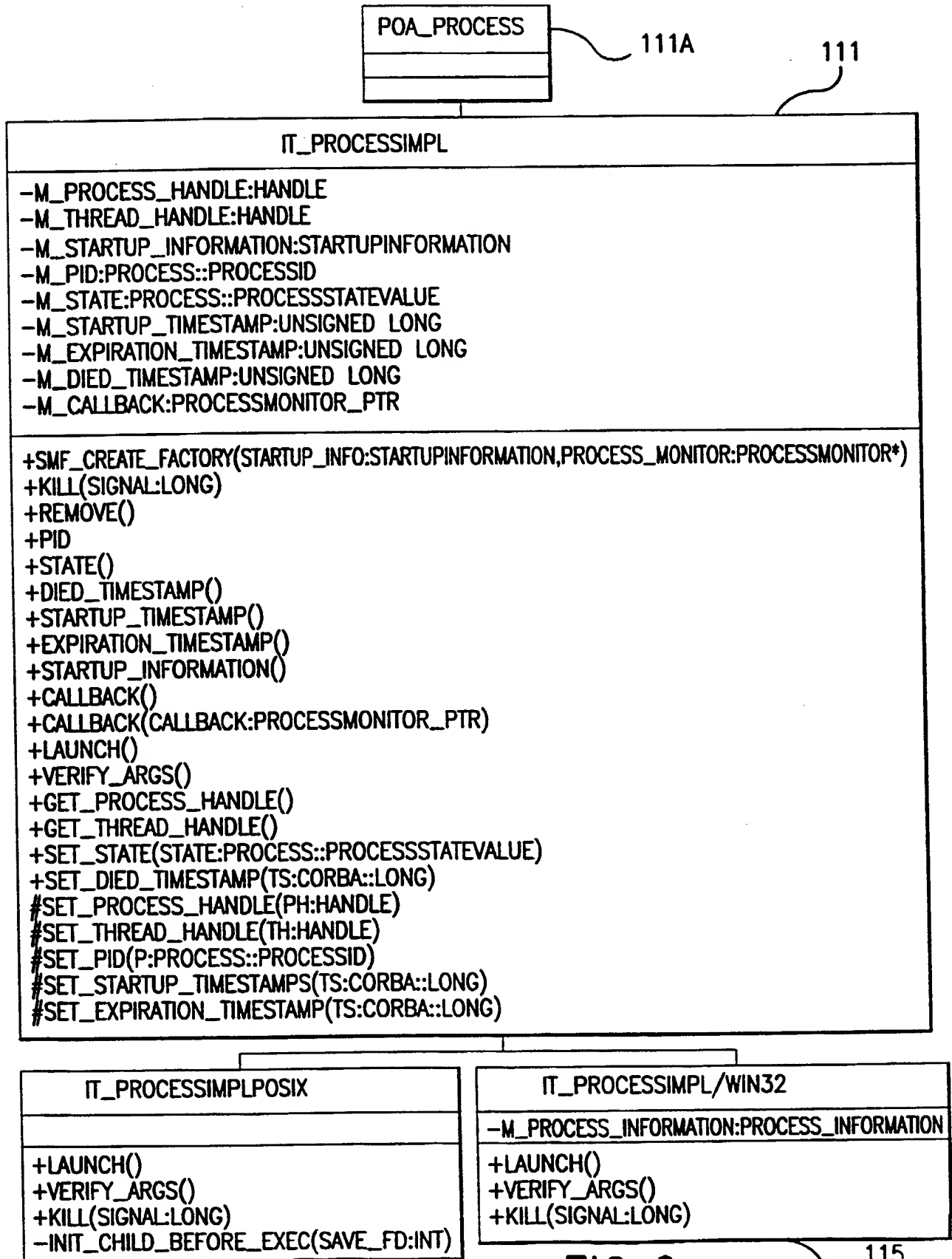


FIG. 9

9/14

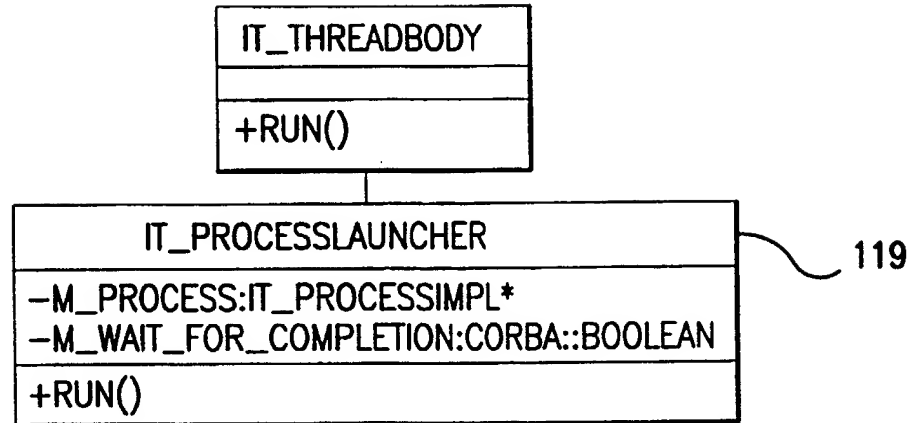


FIG. 10

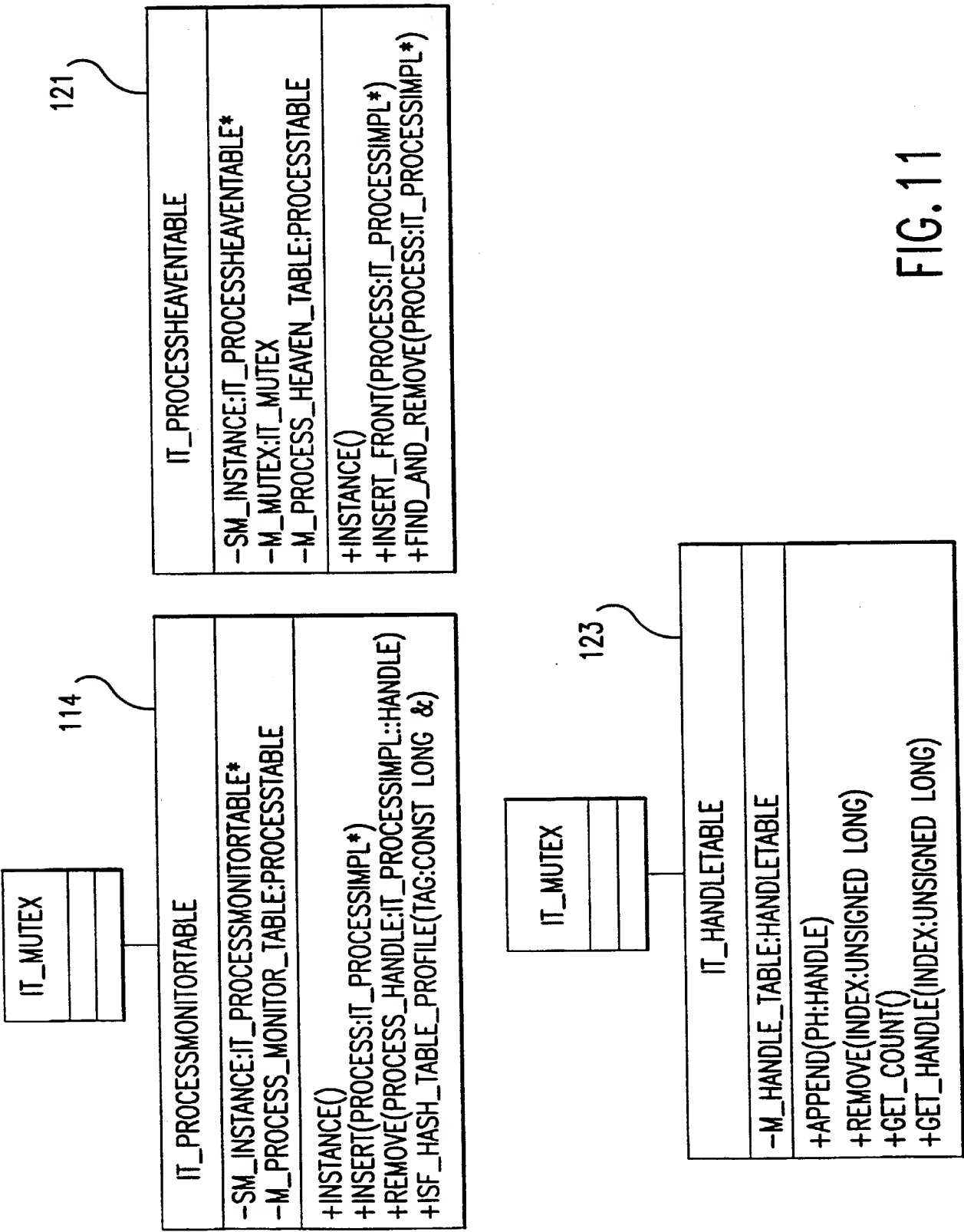


FIG. 11

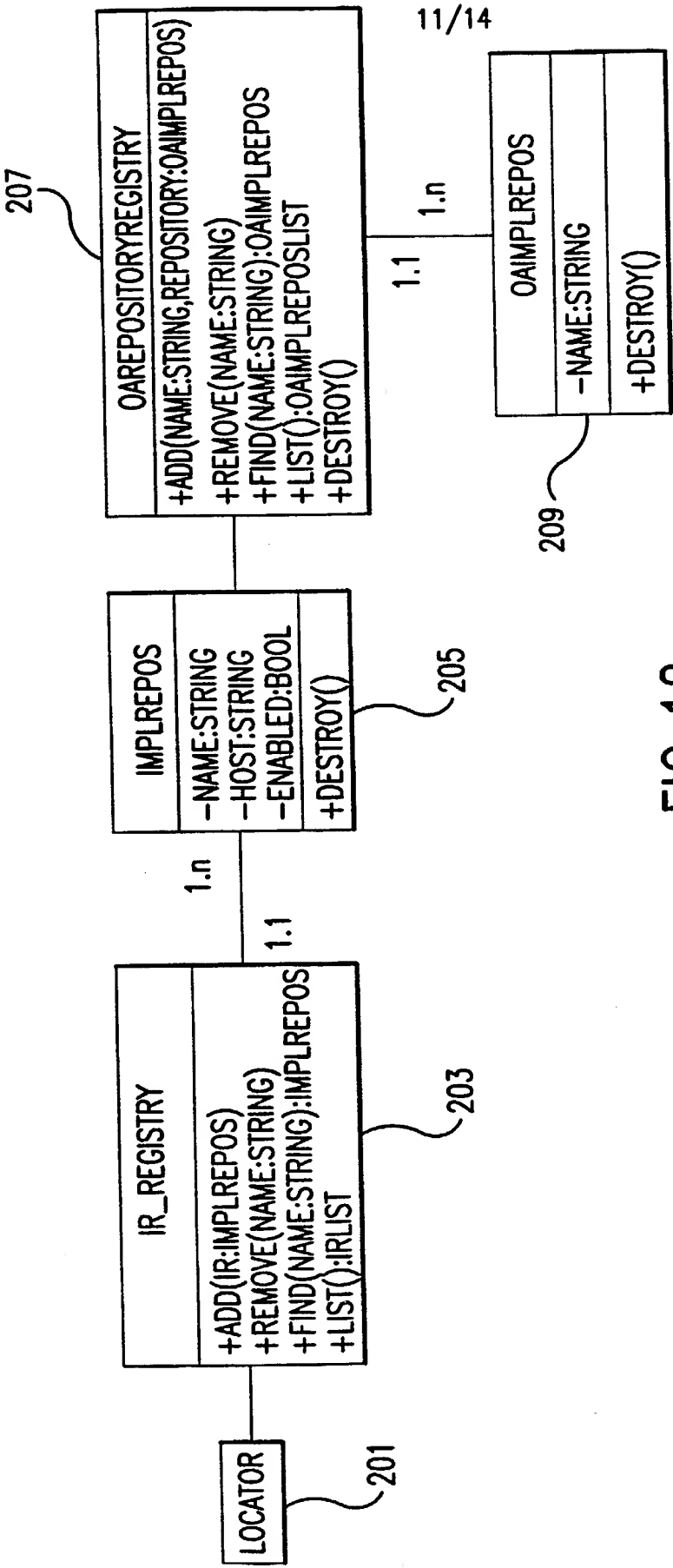


FIG. 12

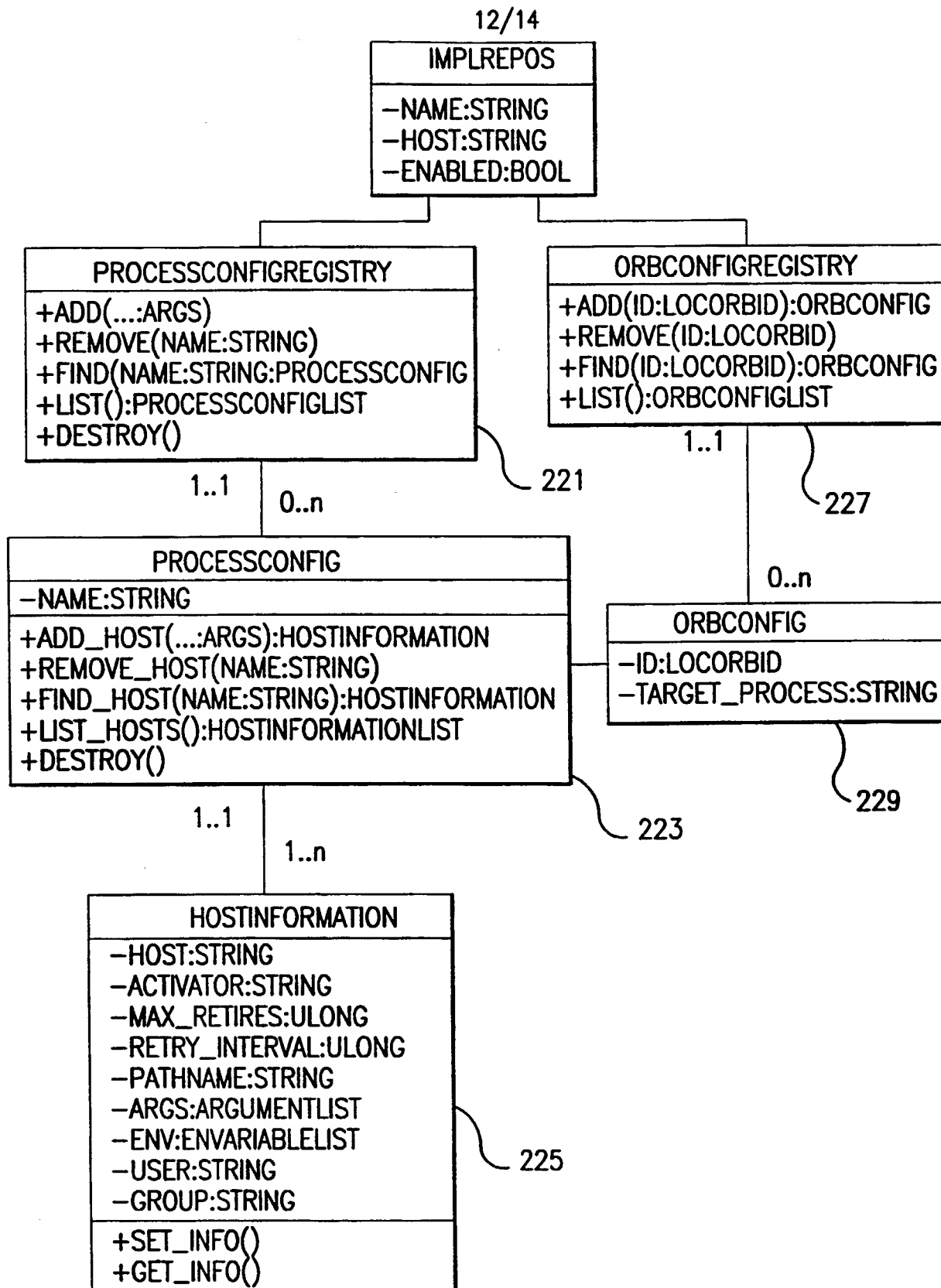


FIG. 13

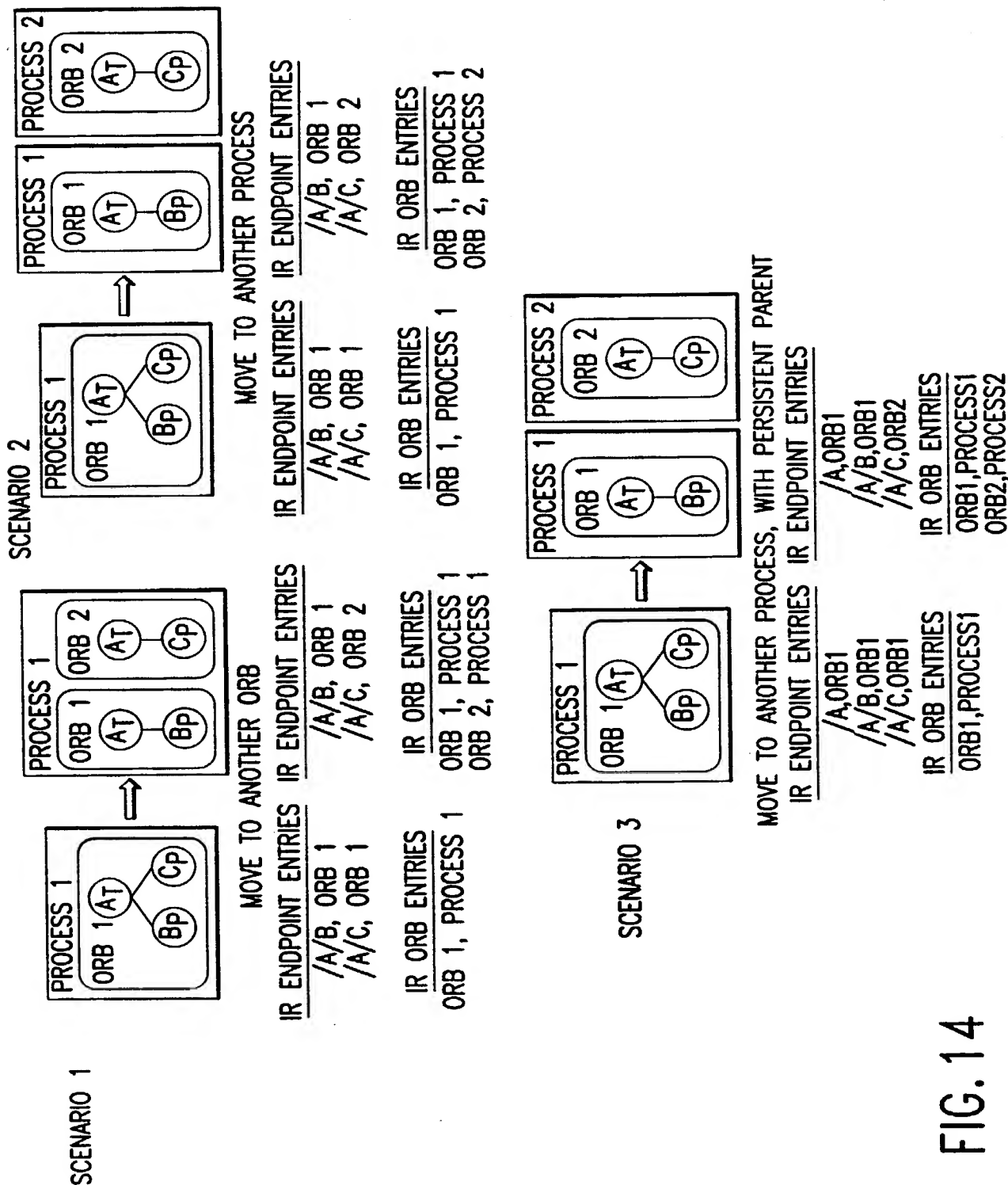


FIG. 14

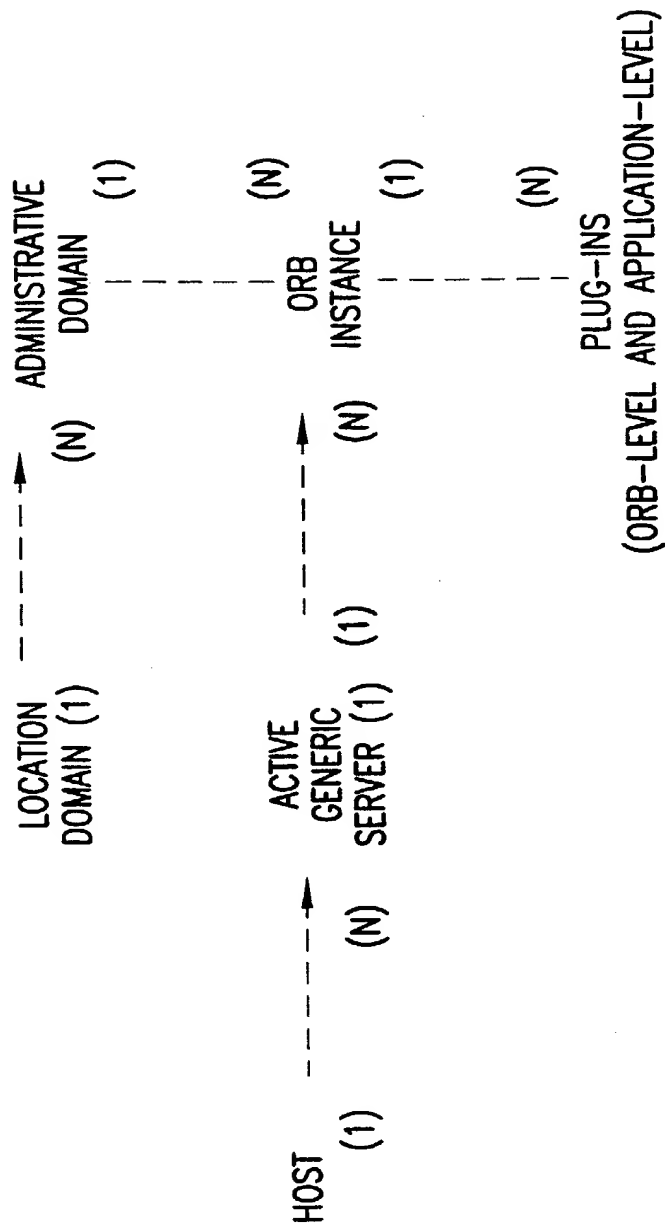


FIG. 15

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 with amendment(s) filed on (if applicable)
☐ was filed as PCT international Application No. on and was amended under PCT Article 19 on (if applicable)

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APPLICATION NUMBER	COUNTRY	DATE OF FILING (day, month, year)	PRIORITY CLAIMED
PCT/US00/02014	PCT	January 28, 2000	YES <input checked="" type="checkbox"/> NO <input type="checkbox"/>
			YES <input type="checkbox"/> NO <input type="checkbox"/>
			YES <input type="checkbox"/> NO <input type="checkbox"/>

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APPLICATION NUMBER	COUNTRY	DATE OF FILING (day, month, year)	PRIORITY CLAIMED
PCT/US00/02014	PCT	January 28, 2000	YES <input checked="" type="checkbox"/> NO <input type="checkbox"/>
			YES <input type="checkbox"/> NO <input type="checkbox"/>
			YES <input type="checkbox"/> NO <input type="checkbox"/>

I hereby claim the benefit under Title 35, United States Code, §119(e) of any United States provisional application(s) listed below.

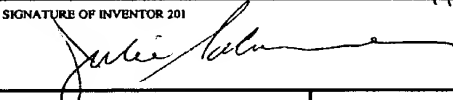
PROVISIONAL APPLICATION NUMBER	FILING DATE

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201	FULL NAME OF INVENTOR	LAST NAME SALAMONE	FIRST NAME Julie	MIDDLE NAME	
	RESIDENCE & CITIZENSHIP	CITY Medford	STATE OR FOREIGN COUNTRY MA MA	COUNTRY OF CITIZENSHIP United States	
	POST OFFICE ADDRESS	STREET 15 Ellis Ave	CITY Medford	STATE OR COUNTRY MA	ZIP CODE 02155
	SIGNATURE OF INVENTOR 201 			DATE 11/21/01	
202	FULL NAME OF INVENTOR	LAST NAME CLARKE	FIRST NAME Alan	MIDDLE NAME	
	RESIDENCE & CITIZENSHIP	CITY	STATE OR FOREIGN COUNTRY	COUNTRY OF CITIZENSHIP United States	
	POST OFFICE ADDRESS	STREET	CITY	STATE OR COUNTRY	ZIP CODE
	SIGNATURE OF INVENTOR 202			DATE	
203	FULL NAME OF INVENTOR	LAST NAME KIELY	FIRST NAME Paul	MIDDLE NAME	
	RESIDENCE & CITIZENSHIP	CITY	STATE OR FOREIGN COUNTRY	COUNTRY OF CITIZENSHIP United States	
	POST OFFICE ADDRESS	STREET	CITY	STATE OR COUNTRY	ZIP CODE
	SIGNATURE OF INVENTOR 203			DATE	
204	FULL NAME OF INVENTOR	LAST NAME WITHAM	FIRST NAME Ronald, Jr.	MIDDLE NAME C	
	RESIDENCE & CITIZENSHIP	CITY	STATE OR FOREIGN COUNTRY	COUNTRY OF CITIZENSHIP United States	
	POST OFFICE ADDRESS	STREET	CITY	STATE OR COUNTRY	ZIP CODE
	SIGNATURE OF INVENTOR 204			DATE	
205	FULL NAME OF INVENTOR	LAST NAME SULLIVAN	FIRST NAME Kevin	MIDDLE NAME	
	RESIDENCE & CITIZENSHIP	CITY	STATE OR FOREIGN COUNTRY	COUNTRY OF CITIZENSHIP United States	
	POST OFFICE ADDRESS	STREET	CITY	STATE OR COUNTRY	ZIP CODE
	SIGNATURE OF INVENTOR 205			DATE	

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	RESIDENCE & CITIZENSHIP	CITY	STATE OR FOREIGN COUNTRY Ireland	COUNTRY OF CITIZENSHIP Ireland	
	POST OFFICE ADDRESS	STREET	CITY	STATE OR COUNTRY Ireland	ZIP CODE
	SIGNATURE OF INVENTOR 201			DATE	
202	FULL NAME OF INVENTOR	LAST NAME CLARKE	FIRST NAME Alan	MIDDLE NAME	
	RESIDENCE & CITIZENSHIP	CITY Dublin 7	STATE OR FOREIGN COUNTRY Ireland IEX	COUNTRY OF CITIZENSHIP Ireland	
	POST OFFICE ADDRESS	STREET Phibsboro, 4 The Borough	CITY Dublin 7	STATE OR COUNTRY Ireland	ZIP CODE
	SIGNATURE OF INVENTOR 202 <i>Alan Clarke</i>			DATE 12th Nov 2002	
203	FULL NAME OF INVENTOR	LAST NAME KIELY	FIRST NAME Paul	MIDDLE NAME	
	RESIDENCE & CITIZENSHIP	CITY Leixlip, Co. Kildare	STATE OR FOREIGN COUNTRY Ireland	COUNTRY OF CITIZENSHIP Ireland	
	POST OFFICE ADDRESS	STREET 8 Rockingham Green, Station Road	CITY Leixlip, Co. Kildare	STATE OR COUNTRY Ireland	ZIP CODE
	SIGNATURE OF INVENTOR 203			DATE	
204	FULL NAME OF INVENTOR	LAST NAME WITHAM	FIRST NAME Ronald, Jr.	MIDDLE NAME C	
	RESIDENCE & CITIZENSHIP	CITY Halibut Bay	STATE OR FOREIGN COUNTRY Canada	COUNTRY OF CITIZENSHIP Ireland	
	POST OFFICE ADDRESS	STREET 308 Ketch Harbour Road	CITY Halibut Bay (NS)	STATE OR COUNTRY Canada	ZIP CODE BeV 1J6
	SIGNATURE OF INVENTOR 204			DATE	
205	FULL NAME OF INVENTOR	LAST NAME SULLIVAN	FIRST NAME Kevin	MIDDLE NAME	
	RESIDENCE & CITIZENSHIP	CITY Cheekpoint, Co. Waterford	STATE OR FOREIGN COUNTRY Ireland	COUNTRY OF CITIZENSHIP Ireland	
	POST OFFICE ADDRESS	STREET Minaun	CITY Cheekpoint, Co. Waterford	STATE OR COUNTRY Ireland	ZIP CODE
	SIGNATURE OF INVENTOR 205			DATE	

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APPLICATION NUMBER	COUNTRY	DATE OF FILING (day, month, year)	PRIORITY CLAIMED
PCT/US00/02014	PCT	January 28, 2000	YES <input checked="" type="checkbox"/> NO <input type="checkbox"/>
			YES <input type="checkbox"/> NO <input type="checkbox"/>
			YES <input type="checkbox"/> NO <input type="checkbox"/>

I hereby claim the benefit under Title 35, United States Code, §119(c) of any United States provisional application(s) listed below.

PROVISIONAL APPLICATION NUMBER	FILING DATE

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PENNIE & EDMONDS LLP DOCKET NO. 9318-022-999

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	POST OFFICE ADDRESS	STREET 15 Ellis Avenue	CITY Medford	STATE OR COUNTRY MA ZIP CODE 02155
SIGNATURE OF INVENTOR 201			DATE	
2 0 2	FULL NAME OF INVENTOR	LAST NAME CLARKE	FIRST NAME Alan	MIDDLE NAME
	RESIDENCE & CITIZENSHIP	CITY Dublin 7	STATE OR FOREIGN COUNTRY Ireland	COUNTRY OF CITIZENSHIP Ireland
	POST OFFICE ADDRESS	STREET Phibsboro, 4 The Borough	CITY Dublin 7	STATE OR COUNTRY Ireland ZIP CODE
SIGNATURE OF INVENTOR 202			DATE	
2 0 3	FULL NAME OF INVENTOR	LAST NAME <u>KIRLY</u>	FIRST NAME Paul	MIDDLE NAME
	RESIDENCE & CITIZENSHIP	CITY Leixlip, Co. Kildare	STATE OR FOREIGN COUNTRY Ireland <u>IE</u>	COUNTRY OF CITIZENSHIP Ireland
	POST OFFICE ADDRESS	STREET 8 Rockingham Green, Station Road	CITY Leixlip, Co. Kildare	STATE OR COUNTRY Ireland ZIP CODE
SIGNATURE OF INVENTOR 203 <u>Paul Kirly</u>			DATE 11 / APRIL / 02	
2 0 4	FULL NAME OF INVENTOR	LAST NAME WITHAM	FIRST NAME Ronald, Jr.	MIDDLE NAME C
	RESIDENCE & CITIZENSHIP	CITY Halibur Bay	STATE OR FOREIGN COUNTRY Canada	COUNTRY OF CITIZENSHIP United States
	POST OFFICE ADDRESS	STREET 308 Ketch Harbour Road	CITY Halibur Bay (NS)	STATE OR COUNTRY Canada ZIP CODE B3V 1J6
SIGNATURE OF INVENTOR 204			DATE	
2 0 5	FULL NAME OF INVENTOR	LAST NAME SULLIVAN	FIRST NAME Kevin	MIDDLE NAME
	RESIDENCE & CITIZENSHIP	CITY Checkpoint, Co. Wexford	STATE OR FOREIGN COUNTRY Ireland	COUNTRY OF CITIZENSHIP Ireland
	POST OFFICE ADDRESS	STREET Mimam	CITY Checkpoint, Co. Wexford	STATE OR COUNTRY Ireland ZIP CODE
SIGNATURE OF INVENTOR 205			DATE	

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PCT/US00/02014	PCT	January 28, 2000	YES <input checked="" type="checkbox"/> NO <input type="checkbox"/>
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	POST OFFICE ADDRESS	STREET 308 Ketch Harbour Road	CITY Halibut Bay (NS)	STATE OR COUNTRY Canada ZIP CODE B3V 1J6
	SIGNATURE OF INVENTOR 204			DATE 4/15/2002
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	POST OFFICE ADDRESS	STREET	CITY	STATE OR COUNTRY Ireland
	SIGNATURE OF INVENTOR 201			DATE
202	FULL NAME OF INVENTOR	LAST NAME CLARKE	FIRST NAME Alan	MIDDLE NAME
	RESIDENCE & CITIZENSHIP	CITY Dublin 7	STATE OR FOREIGN COUNTRY Ireland	COUNTRY OF CITIZENSHIP Ireland
	POST OFFICE ADDRESS	STREET Pinebore, 4 The Borough	CITY Dublin 7	STATE OR COUNTRY Ireland
	SIGNATURE OF INVENTOR 202			DATE
203	FULL NAME OF INVENTOR	LAST NAME KIELY	FIRST NAME Paul	MIDDLE NAME
	RESIDENCE & CITIZENSHIP	CITY Lisdip, Co. Kildare	STATE OR FOREIGN COUNTRY Ireland	COUNTRY OF CITIZENSHIP Ireland
	POST OFFICE ADDRESS	STREET 8 Rockingham Green, Station Road	CITY Lisdip, Co. Kildare	STATE OR COUNTRY Ireland
	SIGNATURE OF INVENTOR 203			DATE
204	FULL NAME OF INVENTOR	LAST NAME WITHAM	FIRST NAME Ronald, Jr.	MIDDLE NAME C
	RESIDENCE & CITIZENSHIP	CITY Halibut Bay	STATE OR FOREIGN COUNTRY Canada	COUNTRY OF CITIZENSHIP Ireland
	POST OFFICE ADDRESS	STREET 308 Ketch Harbour Road	CITY Halibut Bay (NS)	STATE OR COUNTRY Canada
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205	FULL NAME OF INVENTOR	LAST NAME SULLIVAN	FIRST NAME Kevin	MIDDLE NAME
	RESIDENCE & CITIZENSHIP	CITY Cheekpoint, Co. Waterford	STATE OR FOREIGN COUNTRY Ireland	COUNTRY OF CITIZENSHIP Ireland
	POST OFFICE ADDRESS	STREET Minjoun	CITY Cheekpoint, Co. Waterford	STATE OR COUNTRY Ireland
	SIGNATURE OF INVENTOR 205 <i>Kevin Sullivan</i>			DATE 12/4/2002

(2)

NY2-1257729.1

SUPPLEMENTAL AFFIDAVIT

I, Margaret O'Keefe, hereby certify that:

1. Julie Salamone, Alan Clarke, Paul Kiely, Ronald C. Witham Jr., and Kevin Sullivan are the inventors of US patent application 09/890,536, which is a national phase application of PCT/US00/02014 filed on January 28, 2000 and entitled "METHOD AND SYSTEM FOR DYNAMIC CONFIGURATION OF ACTIVATORS IN A CLIENT-SERVER ENVIRONMENT". Each of the above-identified individuals was employed by IONA Technologies at the time the invention of the above-identified application was made.

2. Although IONA Technologies was initially unable to obtain an executed declaration from Alan Clarke, it has now succeeded in doing so.

3. Each declaration of the four previously submitted declarations executed by Julie Salamone, Paul Kiely, Ronald C. Witham Jr., and Kevin Sullivan, respectively, and the afore-mentioned declaration executed by Alan Clarke identifies the correct name of all five joint inventors, i.e., Julie Salamone, Alan Clarke, Paul Kiely, Ronald C. Witham Jr., and Kevin Sullivan.

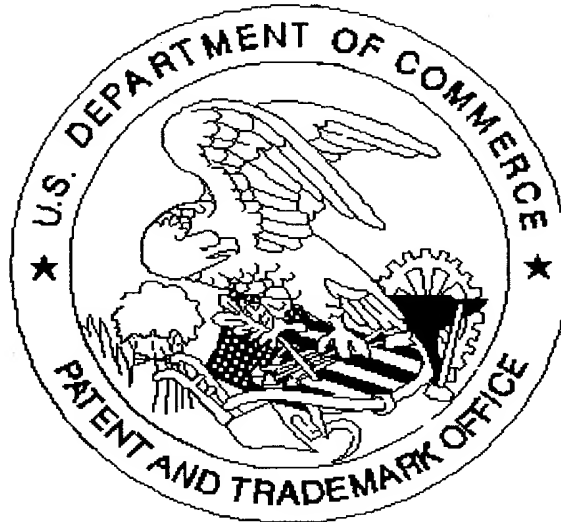
4. The inventor information for each inventor is correct on the copy of the declaration executed by that inventor.

5. I declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed true; and further that these statements were made with knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under 18 U.S.C. § 1001, and that such willful false statements may jeopardize the validity of the application or any patent issuing thereon.

Date: 7th January 2003

Margaret O'Keefe
Margaret O'Keefe
Executive Assistant - SVP Engineering
IONA Technologies

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declaration have line.*